

Principles of Programming Languages

Lecture 1 – Overview and Introduction to OCaml

Ziyang Li

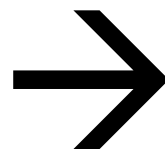
Instructor



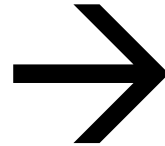
Ziyang Li

- Assistant professor @JHU CS, 2025-
- Before that: PhD at University of Pennsylvania
- Research areas: PL + ML + Sec
- Favorite PL: Rust & JavaScript

Programming...



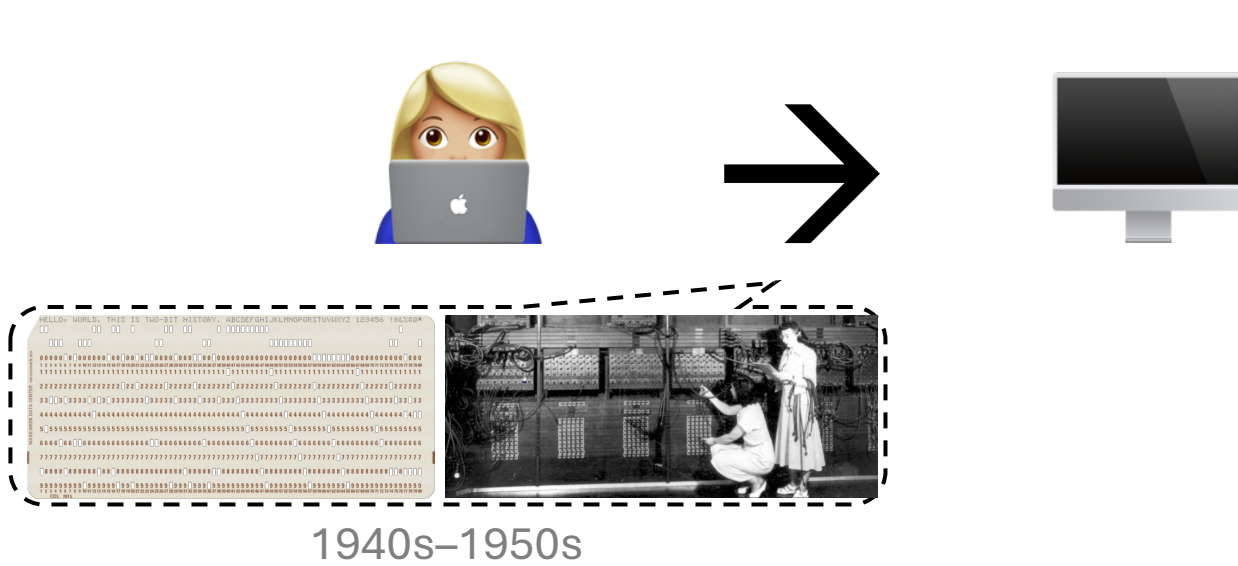
Programming...

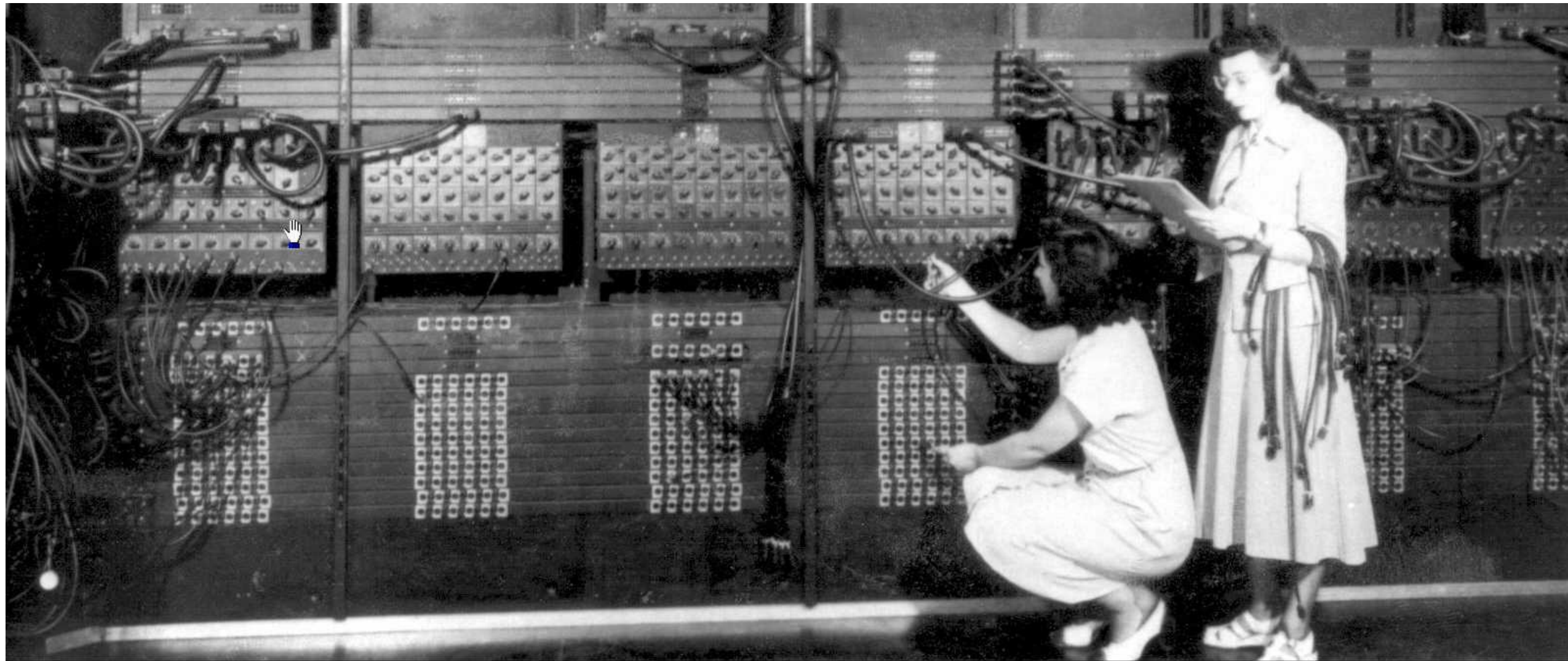


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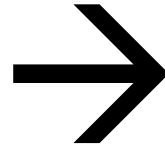
1940s–1950s: Binary

Programming...





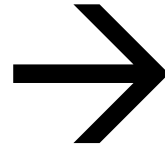
Programming...



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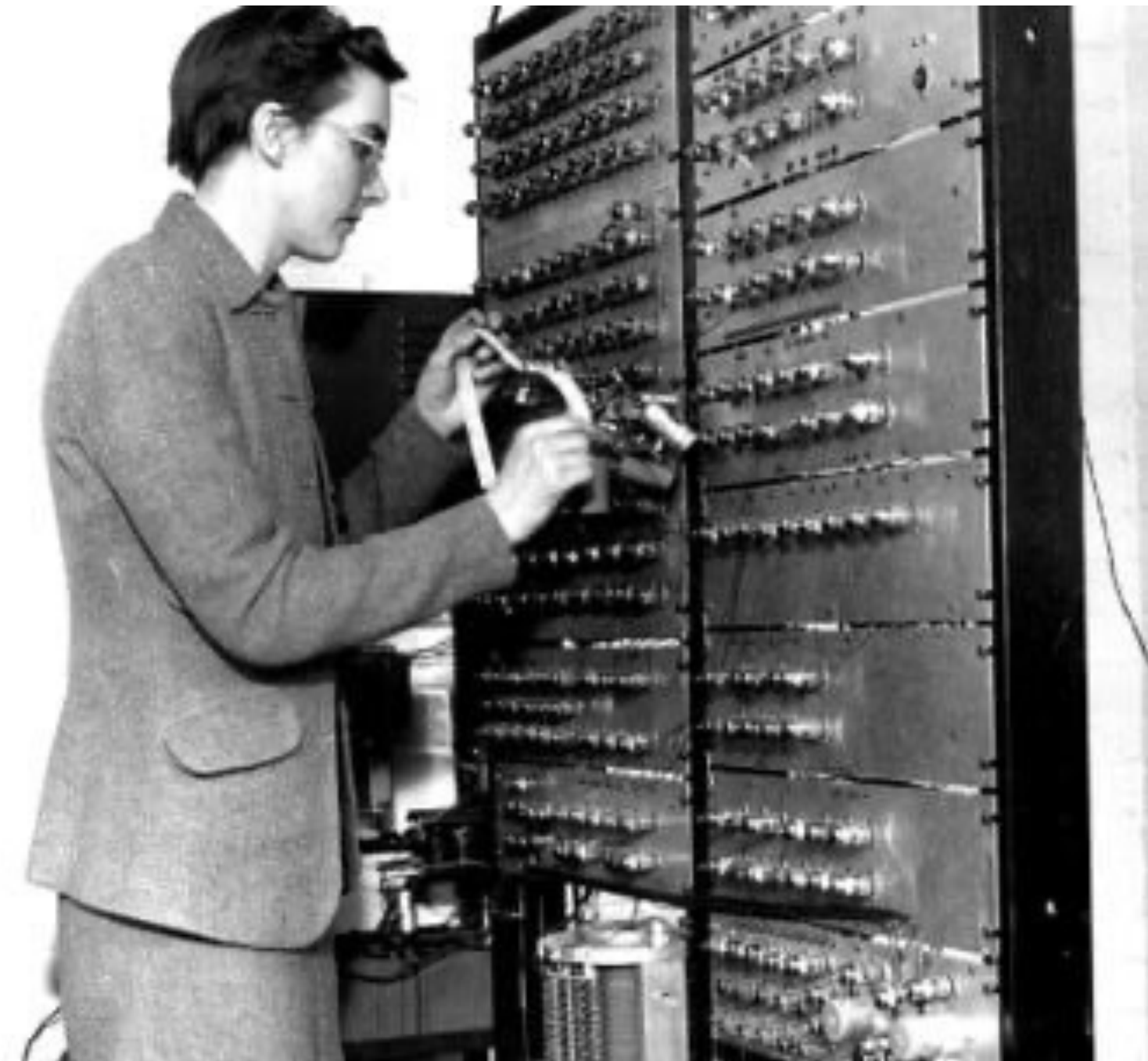
1940s–1950s: Binary

Programming...



ADD R1, R2

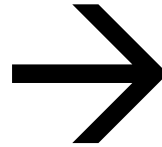
1950s–1960s: Assembly



Kathleen Booth

1922–2022

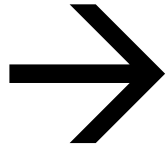
Programming...



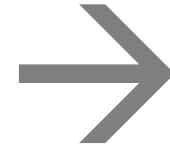
ADD R1, R2

1950s–1960s: Assembly

Programming...



Assembler



ADD R1, R2

1950s–1960s: Assembly

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The FORTRAN Automatic Coding System

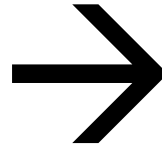
J. W. BACKUS†, R. J. BEEBER†, S. BEST‡, R. GOLDBERG†, L. M. HAIBT†,
H. L. HERRICK†, R. A. NELSON†, D. SAYRE†, P. B. SHERIDAN†,
H. STERN†, I. ZILLER†, R. A. HUGHES§, AND R. NUTT||

INTRODUCTION

THE FORTRAN project was begun in the summer of 1954. Its purpose was to reduce by a large factor the task of preparing scientific problems for IBM's next large computer, the 704. If it were possible for the 704 to code problems for itself and produce as

system is now complete. It has two components: the FORTRAN language, in which programs are written, and the translator or executive routine for the 704 which effects the translation of FORTRAN language programs into 704 programs. Descriptions of the FORTRAN language and the translator form the principal

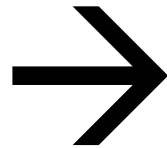
Programming...



```
for (i = 0; i < n; i++)  
    sum += a[i];
```

1970s: High-level Language

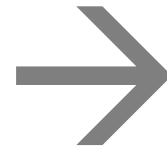
Programming...



Compiler



Assembler



```
for (i = 0; i < n; i++)  
    sum += a[i];
```

1970s: High-level Language

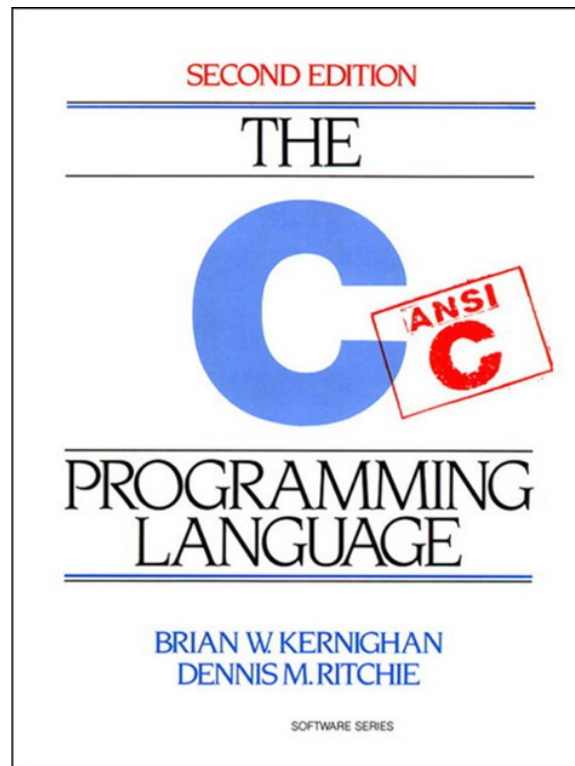
```
ADD R1, R2
```

```
0101010101
```

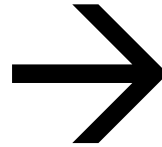


Dennis M. Ritchie

1941–2011



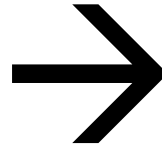
Programming...



```
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1970s: High-level Language

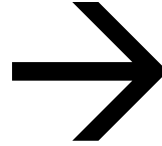
Programming...



```
for i in range(n):  
    sum += a[i]
```

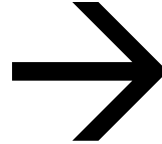
1980s onwards: Modern High-Level Languages

Programming...



```
for i in range(n):  
    sum += a[i]
```

Programming...

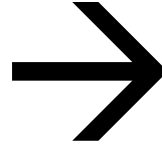


```
for i in range(n):  
    sum += a[i]
```



```
for (var e of a) {  
    sum += e; }  
}
```


Programming...



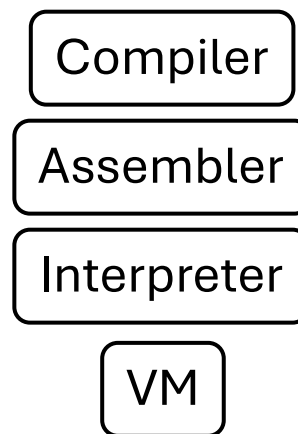
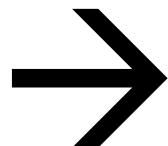
```
for i in range(n):  
    sum += a[i]
```



```
for (var e of a) {  
    sum += e; }  
}
```



```
total :: Num a => [a] -> a  
total a = foldl (+) 0 a
```



...



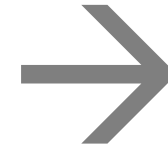
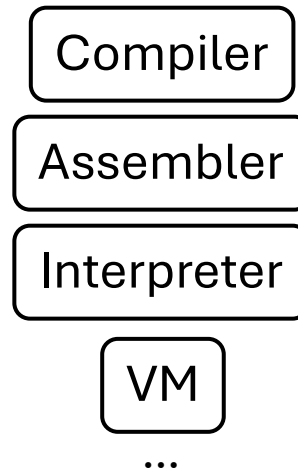
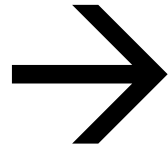
```
for i in range(n):  
    sum += a[i]
```



```
for (var e of a) {  
    sum += e; }  
}
```



```
total :: Num a => [a] -> a  
total a = foldl (+) 0 a
```



```
for i in range(n):  
    sum += a[i]
```



```
for (var e of a) {  
    sum += e; }  
}
```

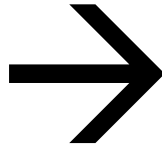


```
total :: Num a => [a] -> a  
total a = foldl (+) 0 a
```

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ADD R1, R2

Programming...



```
for i in range(n):  
    sum += a[i]
```

Assembly

```
append:
    push ebp
    mov ebp, esp
    push eax
    push ebx
    push len
    call malloc
    mov ebx, [ebp + 12]
    mov [eax + info], ebx
    mov dword [eax + next], 0
    mov ebx, [ebp + 8]
    cmp dword [ebx], 0
    je null_pointer
    mov ebx, [ebx]

next_element:
    cmp dword [ebx + next], 0
    je found_last
    mov ebx, [ebx + next]
    jmp next_element

found_last:
    push eax
    push addMes
    call puts
    add esp, 4
    pop eax
    mov [ebx + next], eax

go_out:
    pop ebx
    pop eax
    mov esp, ebp
    pop ebp
    ret 8

null_pointer:
    push eax
    push nullMes
    call puts
    add esp, 4
    pop eax
    mov [ebx], eax
    jmp go_out
```

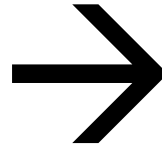
C

```
void insert(node *xs, int x) {
    node *new;
    node *temp;
    node *prev;
    new = (node *)malloc(sizeof(node));
    if(new == NULL) {
        printf("Insufficient memory.");
        return;
    }
    new->val = x;
    new->next = NULL;
    if (xs == NULL) {
        xs = new;
    } else if(x < xs->val) {
        new->next = xs;
        xs = new;
    } else {
        prev = xs;
        temp = xs->next;
        while(temp != NULL && x > temp->val) {
            prev = temp;
            temp = temp->next;
        }
        if(temp == NULL) {
            prev->next = new;
        } else {
            new->next = temp;
            prev->next = new;
        }
    }
}
```

OCaml

```
let rec insert x ys =
  match ys with
  | [] -> [x]
  | y :: rs ->
    if x <= y then x :: y :: rs
    else y :: (insert x rs)
```

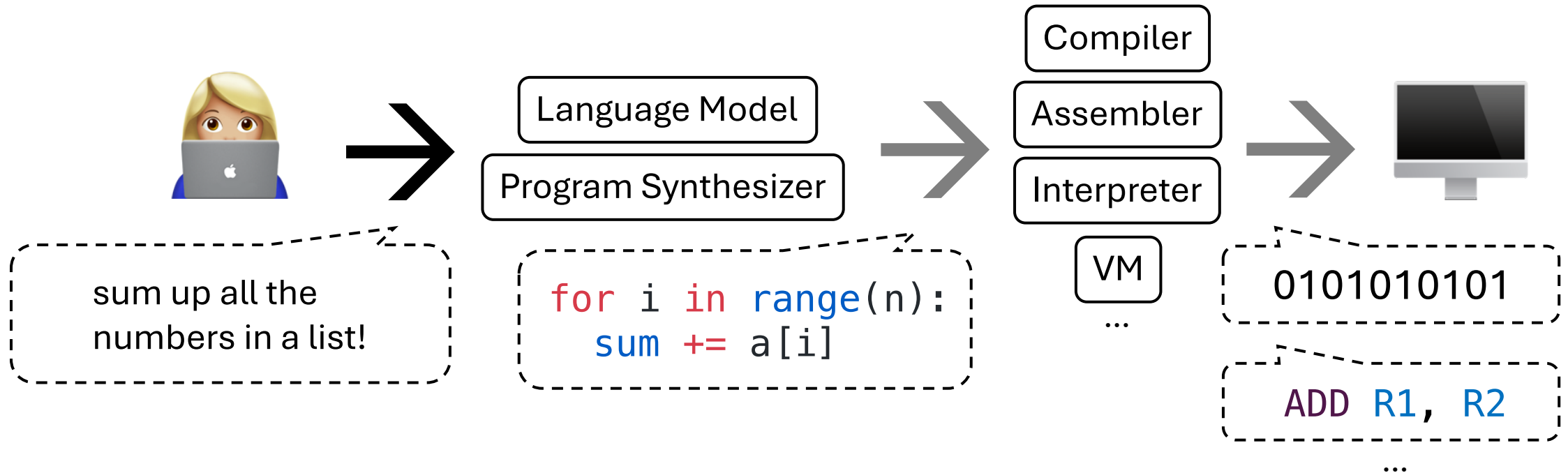
Programming → Talking to Computers

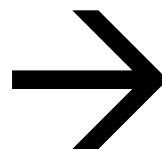


sum up all the numbers in a list!

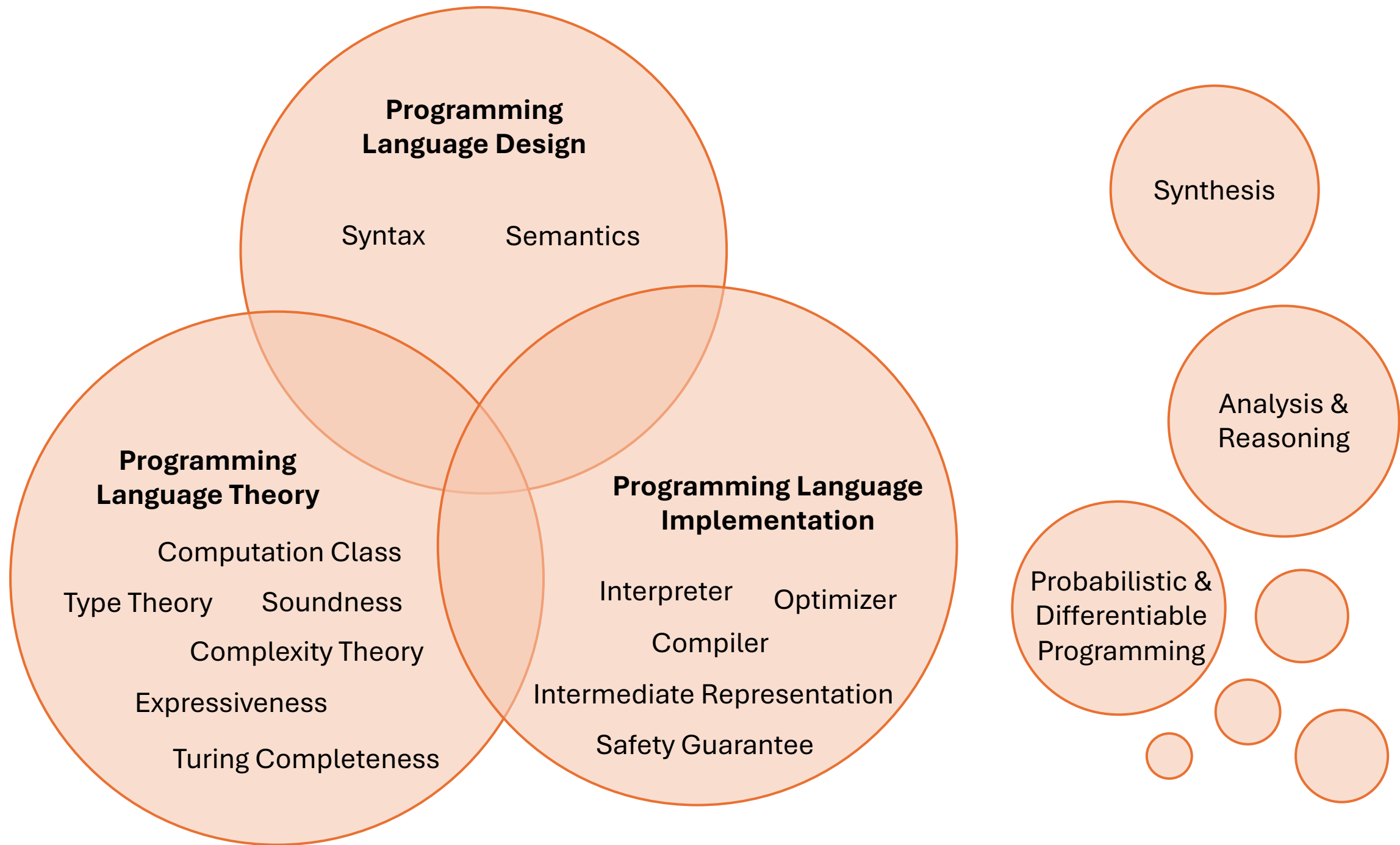
Natural language as programming languages?

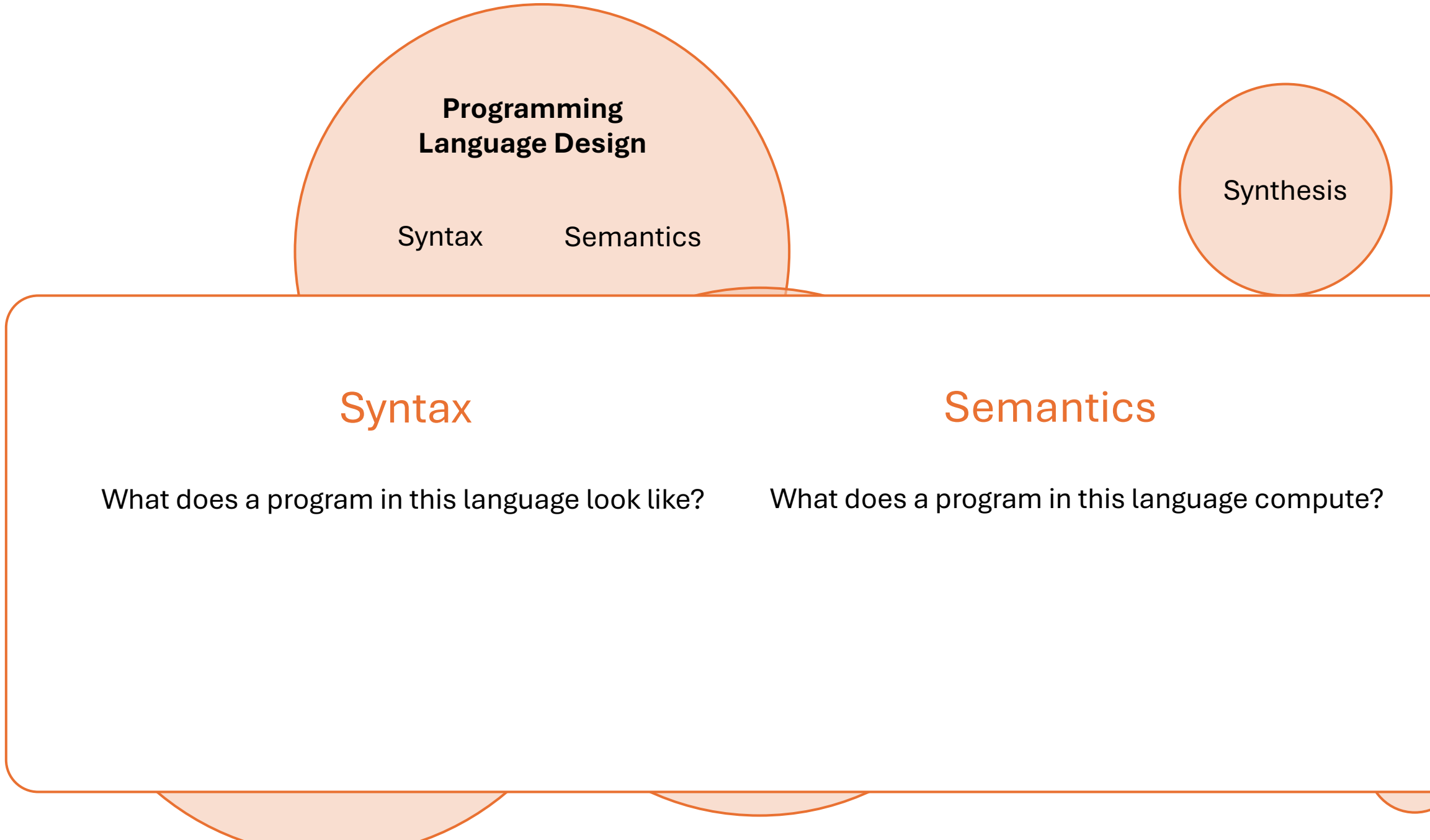
Programming → Talking to Computers





Principles of Programming Languages





The diagram illustrates the components of Programming Language Design. At the top, a large light-orange semi-circle contains the text 'Programming Language Design' and 'Syntax Semantics'. To its right is a smaller light-orange circle containing the text 'Synthesis'. Below these, a large light-orange rounded rectangle contains the text 'Syntax' and 'Semantics' in orange, followed by their respective questions: 'What does a program in this language look like?' and 'What does a program in this language compute?'. The entire diagram is framed by a light-orange border with semi-circular ends at the top and bottom.

Programming Language Design

Syntax

Semantics

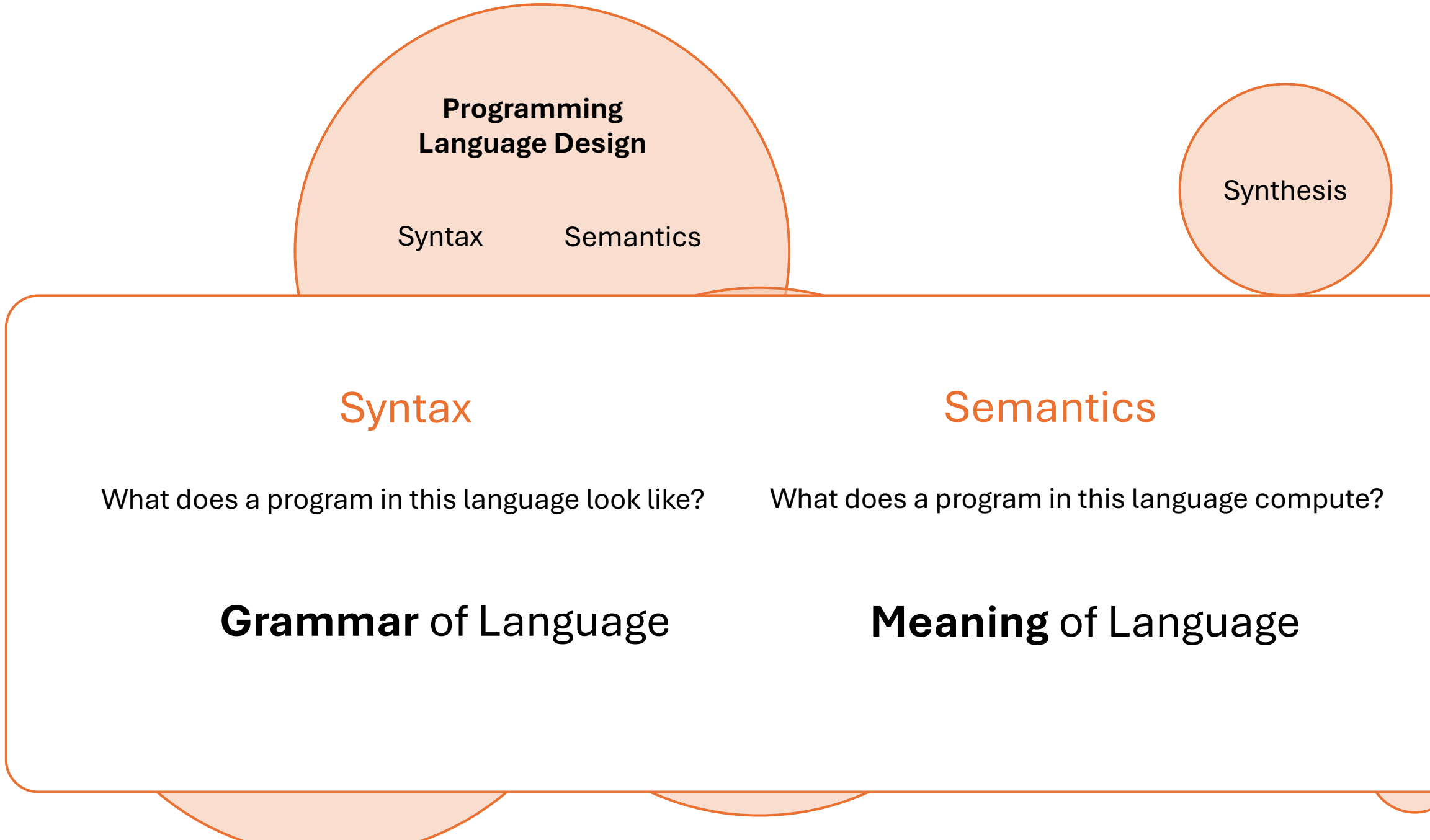
Synthesis

Syntax

What does a program in this language look like?

Semantics

What does a program in this language compute?



The diagram illustrates the components of Programming Language Design. At the top, a large light-orange semi-circle contains the text 'Programming Language Design' and 'Syntax Semantics'. To its right is a smaller light-orange circle containing the text 'Synthesis'. Below these, a large white rounded rectangle with an orange border contains the text 'Syntax' and 'Semantics' in orange, followed by their respective questions: 'What does a program in this language look like?' and 'What does a program in this language compute?'. At the bottom of this rectangle are the bolded terms 'Grammar of Language' and 'Meaning of Language'. The entire diagram is framed by light-orange semi-circular shapes at the top and bottom.

Programming Language Design

Syntax

Semantics

Synthesis

Syntax

Semantics

What does a program in this language look like?

What does a program in this language compute?

Grammar of Language

Meaning of Language


```

if_stmt:
    | 'if' named_expression ':' block elif_stmt
    | 'if' named_expression ':' block [else_block]
elif_stmt:
    | 'elif' named_expression ':' block elif_stmt
    | 'elif' named_expression ':' block [else_block]
else_block:
    | 'else' ':' block

# While statement
# -----

while_stmt:
    | 'while' named_expression ':' block [else_block]

# For statement
# -----

for_stmt:
    | 'for' star_targets 'in' ~ star_expressions ':' [TYPE_COMMENT] block [else_block]
    | 'async' 'for' star_targets 'in' ~ star_expressions ':' [TYPE_COMMENT] block [el

# With statement
# -----

with_stmt:
    | 'with' '(' ','.with_item+ ','? ')' ':' [TYPE_COMMENT] block
    | 'with' ','.with_item+ ':' [TYPE_COMMENT] block
    | 'async' 'with' '(' ','.with_item+ ','? ')' ':' block
    | 'async' 'with' ','.with_item+ ':' [TYPE_COMMENT] block

with_item:
    | expression 'as' star_target &(',' | ')') | ':'
    | expression

# Try statement
# -----

try_stmt:
    | 'try' ':' block finally_block
    | 'try' ':' block except_block+ [else_block] [finally_block]
    | 'try' ':' block except_star_block+ [else_block] [finally_block]

```

What language is this?

Concrete and Abstract Syntax

Python Grammar (Concrete Syntax)

```
if_stmt:
    'if' named_expression ':' block elif_stmt
    'if' named_expression ':' block [else_block]
elif_stmt:
    'elif' named_expression ':' block elif_stmt
    'elif' named_expression ':' block [else_block]
else_block:
    'else' ':' block

# While statement
# -----
while_stmt:
    'while' named_expression ':' block [else_block]

# For statement
# -----
for_stmt:
    'for' star_targets 'in' ~ star_expressions ':' [TYPE_COMMENT] block [else_block]
    'async' 'for' star_targets 'in' ~ star_expressions ':' [TYPE_COMMENT] block [el

# With statement
# -----
with_stmt:
    'with' '(' ',' with_item+ ',' '?' ')' ':' [TYPE_COMMENT] block
    'with' ',' with_item+ ':' [TYPE_COMMENT] block
    'async' 'with' '(' ',' with_item+ ',' '?' ')' ':' block
    'async' 'with' ',' with_item+ ':' [TYPE_COMMENT] block

with_item:
    expression 'as' star_target &(',' | '|') | ':'
    expression

# Try statement
# -----
try_stmt:
    'try' ':' block finally_block
    'try' ':' block except_block+ [else_block] [finally_block]
    'try' ':' block except_star_block+ [else_block] [finally_block]
```

<https://docs.python.org/3/reference/grammar.html>

Python Abstract Syntax

```
-- use 'orelse' because else is a keyword in target languages
For(expr target, expr iter, stmt* body, stmt* orelse, string?
AsyncFor(expr target, expr iter, stmt* body, stmt* orelse, st
While(expr test, stmt* body, stmt* orelse)
If(expr test, stmt* body, stmt* orelse)
With(withitem* items, stmt* body, string? type_comment)
AsyncWith(withitem* items, stmt* body, string? type_comment)

Match(expr subject, match_case* cases)

Raise(expr? exc, expr? cause)
Try(stmt* body, excepthandler* handlers, stmt* orelse, stmt*
TryStar(stmt* body, excepthandler* handlers, stmt* orelse, st
Assert(expr test, expr? msg)

Import(alias* names)
ImportFrom(identifier? module, alias* names, int? level)

Global(identifier* names)
Nonlocal(identifier* names)
Expr(expr value)
Pass | Break | Continue

-- col_offset is the byte offset in the utf8 string the parser
attributes (int lineno, int col_offset, int? end_lineno, int? e

-- BoolOp() can use left & right?
expr = BoolOp(boolop op, expr* values)
NamedExpr(expr target, expr value)
BinOp(expr left, operator op, expr right)
UnaryOp(unaryop op, expr operand)
Lambda(arguments args, expr body)
IfExp(expr test, expr body, expr orelse)
Dict(expr* keys, expr* values)
Set(expr* elts)
ListComp(expr elt, comprehension* generators)
```

<https://docs.python.org/3/library/ast.html>

Formal Semantics

WebAssembly (WASM)

Datafun

Reduction

$$\begin{array}{c}
 \frac{s; v^*; e^* \hookrightarrow_i s'; v'^*; e'^*}{s; v^*; L^k[e^*] \hookrightarrow_i s'; v'^*; L^k[e'^*]} \quad \frac{s; v^*; e^* \hookrightarrow_i s'; v'^*; e'^*}{s; v_0^*; \text{local}_n\{i; v^*\} e^* \text{ end} \hookrightarrow_j s'; v_0^*; \text{local}_n\{i; v'^*\} e'^* \text{ end}} \quad \boxed{s; v^*; e^* \hookrightarrow_i s; v^*; e^*} \\
 \\
 L^0[\text{trap}] \hookrightarrow \text{trap} \quad \text{if } L^0 \neq [-] \\
 \frac{(t.\text{const } c) t.\text{unop}}{(t.\text{const } c_1) (t.\text{const } c_2) t.\text{binop}} \hookrightarrow t.\text{const } \text{unop}_t(c) \\
 \frac{(t.\text{const } c_1) (t.\text{const } c_2) t.\text{binop}}{(t.\text{const } c_1) (t.\text{const } c_2) t.\text{binop}} \hookrightarrow t.\text{const } c \quad \text{if } c = \text{binop}_t(c_1, c_2) \\
 \frac{(t.\text{const } c_1) (t.\text{const } c_2) t.\text{binop}}{(t.\text{const } c_1) (t.\text{const } c_2) t.\text{binop}} \hookrightarrow \text{trap} \quad \text{otherwise}
 \end{array}$$

Expression semantics

$$\begin{array}{c}
 \boxed{\alpha : \mathcal{U} \rightarrow \mathcal{U}, \quad \beta : \mathcal{U} \rightarrow \text{Bool}, \quad g : \mathcal{U} \rightarrow \mathcal{U}, \quad \llbracket e \rrbracket : \mathcal{F}_T \rightarrow \mathcal{U}_T} \\
 \\
 \frac{(t.\text{const } c) \quad t :: p(u) \in F_T}{(t.\text{const } c) \quad t :: p(u) \in F_T} \text{ (PREDICATE)} \quad \frac{t :: u \in \llbracket e \rrbracket(F_T) \quad \beta(u) = \text{true}}{t :: u \in \llbracket e \rrbracket(F_T)} \text{ (SELECT)} \quad \frac{t :: u \in \llbracket e \rrbracket(F_T) \quad u' = \alpha(u)}{t :: u' \in \llbracket \pi_\alpha(e) \rrbracket(F_T)} \text{ (PROJECT)} \\
 \\
 \frac{t :: u_2 \in \llbracket e_2 \rrbracket(F_T)}{e_1 \times e_2(F_T)} \text{ (PRODUCT)} \quad \frac{t_2 :: u \in \llbracket e_2 \rrbracket(F_T)}{e_1 - e_2(F_T)} \text{ (DIFF-2)} \quad \frac{u \in g(\{u_i\}_{i=1}^n)}{u \in g(\{u_i\}_{i=1}^n)} \text{ (AGGREGATE)}
 \end{array}$$

SQL

$$\left[\begin{array}{l} \text{SELECT DISTINCT } \alpha : \beta' \mid * \\ \text{FROM } \tau : \beta \text{ WHERE } \theta \end{array} \right]_{D, \eta, x} = \varepsilon \left(\left[\begin{array}{l} \text{SELECT } \alpha : \beta' \mid * \\ \text{FROM } \tau : \beta \text{ WHERE } \theta \end{array} \right]_{D, \eta, x} \right)$$

Signature

$\text{alloc}(\tau_1, \dots, \tau_n)(\overline{r_n}, S)$
 $\overline{d_m} \leftarrow \text{eval}(\alpha_{n,m})(\overline{s_n})$
 $\overline{d_n} \leftarrow \text{gather}(i, \overline{s_n})$
 $d \leftarrow \text{gather}(\alpha_{n,1})(\overline{i_n}, \overline{s_n})$
 $\text{store}(\rho)(\overline{s_n}, s_r)$
 $[\overline{s_n}, s_r] = \text{load}(\rho)()$
 $\overline{d_n} \leftarrow \text{build}(\overline{s_n})$
 $\overline{d_n} \leftarrow \text{count}(\overline{b_n}, h, \overline{a_n})$
 $\overline{d_n} \leftarrow \text{scan}(s)$
 $[d_l, d_r] \leftarrow \text{join}(W)(\overline{b_m}, \overline{d_n})$
 $\overline{d_n} \leftarrow \text{copy}(\overline{s_n})$
 $\overline{d_n} \leftarrow \text{sort}(\overline{s_n})$
 $[\overline{d_n}, s] \leftarrow \text{unique}(\sigma)(\overline{s_n})$
 $\overline{d_n} \leftarrow \text{merge}(\overline{a_n}, \overline{b_n})$

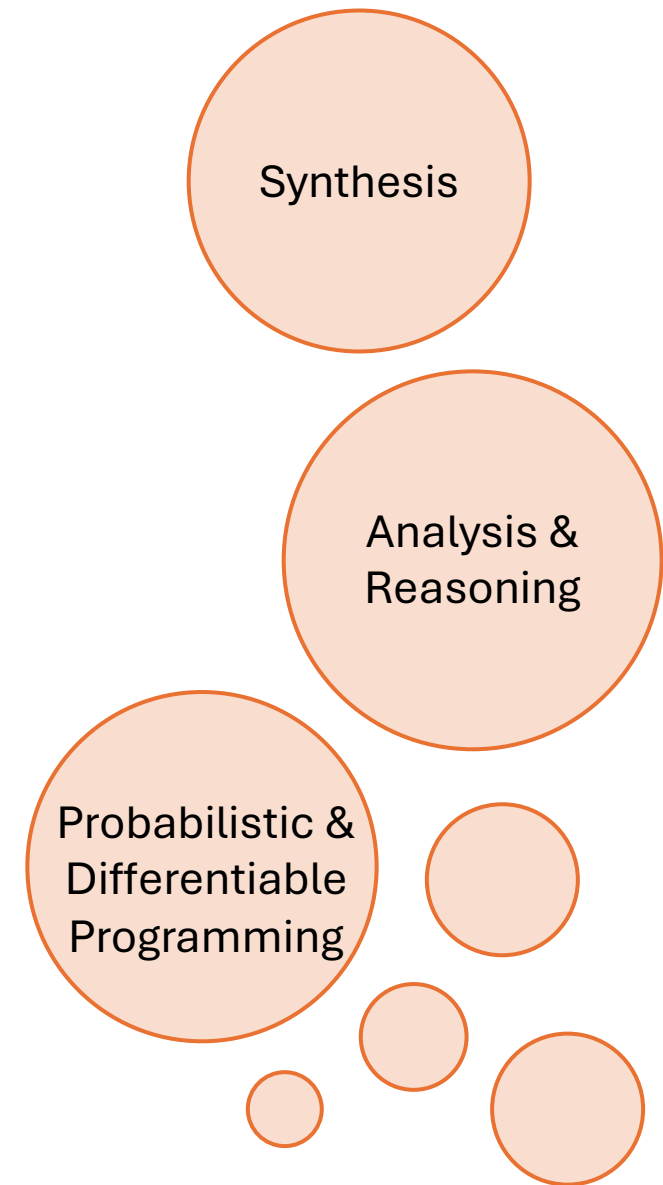
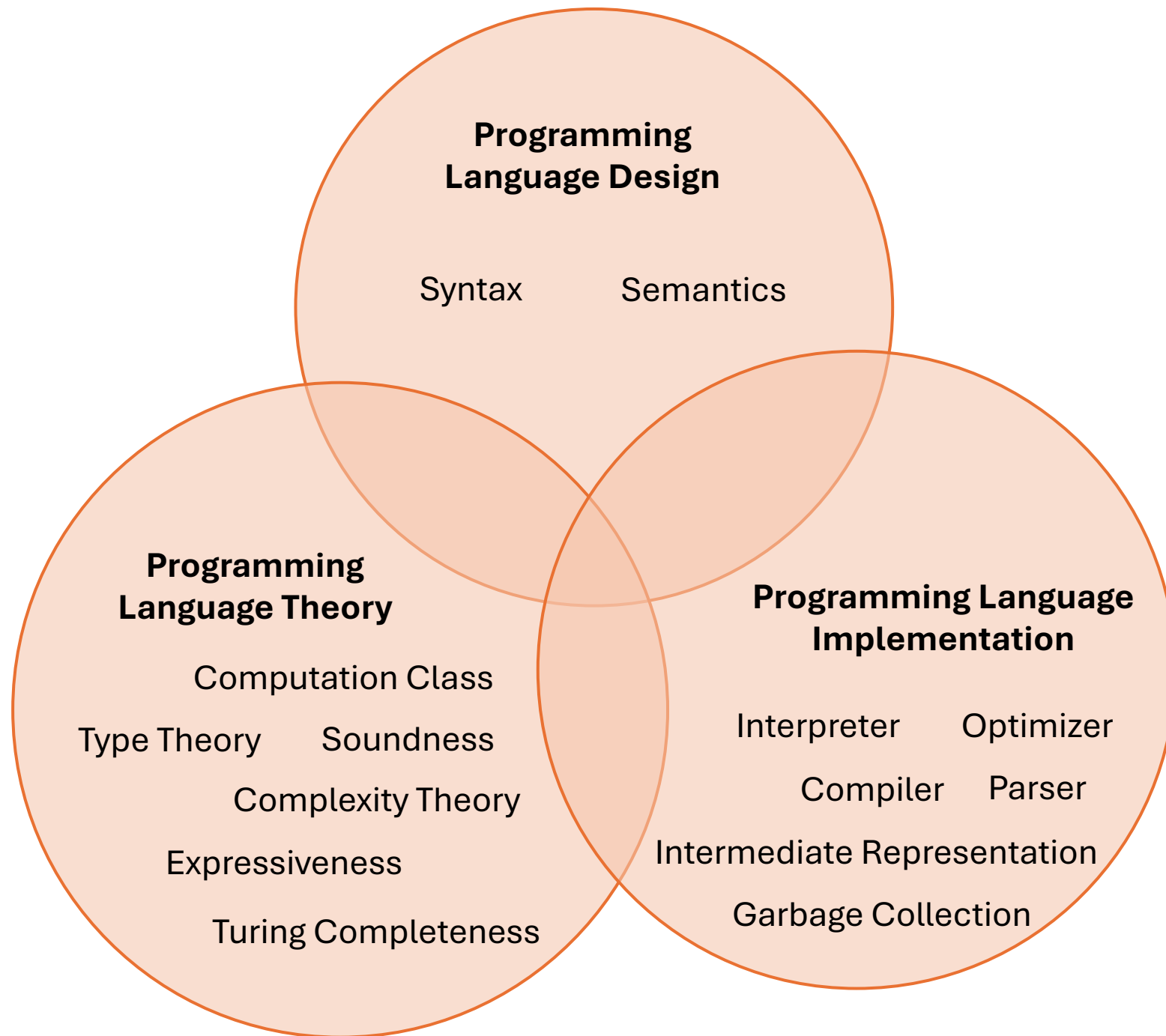
$\llbracket A \rrbracket \in \text{Poset}_0$
 $\llbracket 2 \rrbracket = 2$
 $\llbracket \mathbb{N} \rrbracket = \mathbb{N}_{\leq}$
 $\llbracket \text{str} \rrbracket = \text{Disc } S$
 $\llbracket A \times B \rrbracket = \llbracket A \rrbracket \times \llbracket B \rrbracket$
 $\llbracket A + B \rrbracket = \llbracket A \rrbracket + \llbracket B \rrbracket$
 $\llbracket A \xrightarrow{+} B \rrbracket = \llbracket A \rrbracket \Rightarrow \llbracket B \rrbracket$
 $\llbracket A \rightarrow B \rrbracket = \text{Disc } |\llbracket A \rrbracket| \Rightarrow \llbracket B \rrbracket$
 $\llbracket \{A\} \rrbracket = \mathcal{P}_{\text{fin}} |\llbracket A \rrbracket|$

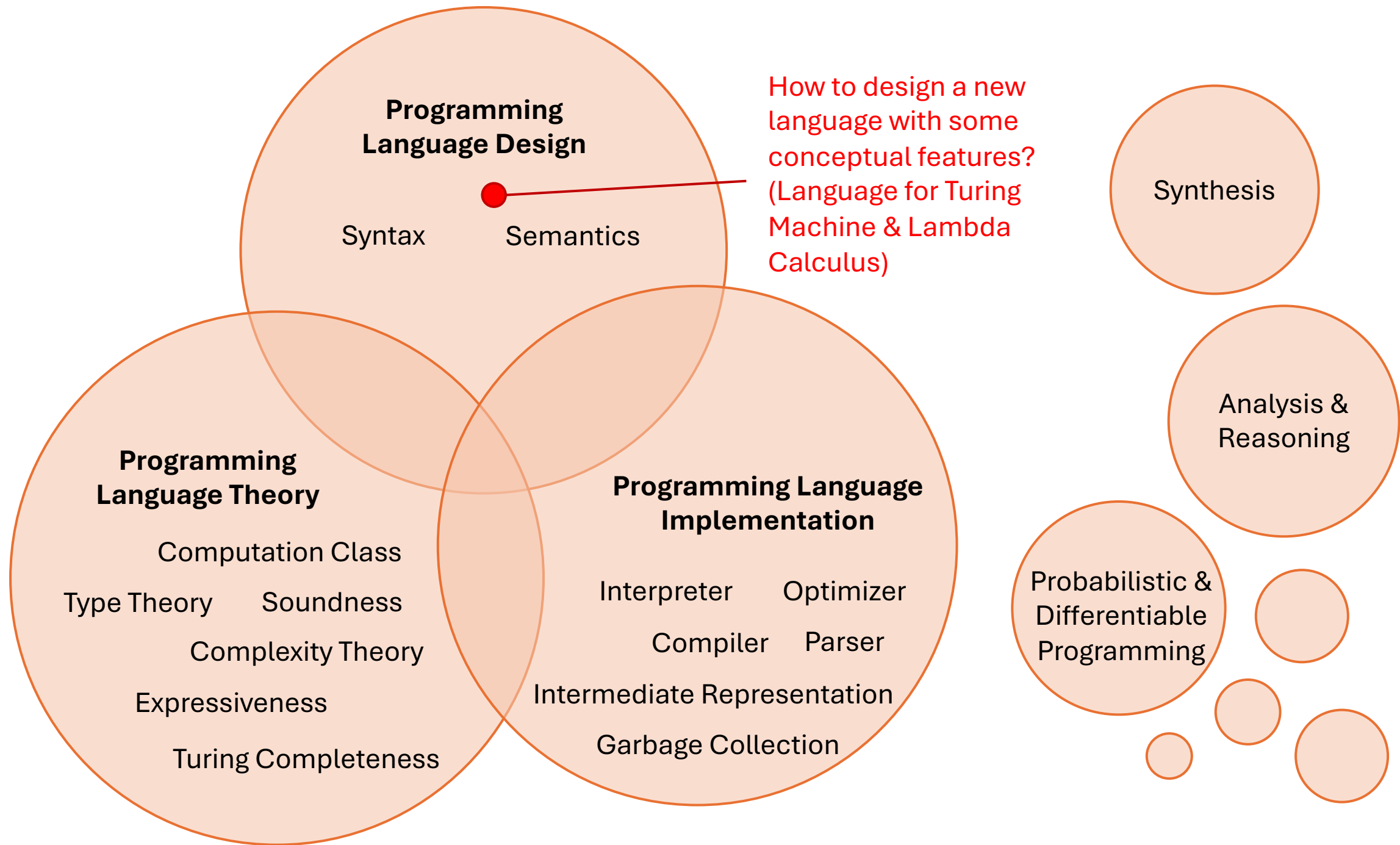
 $\llbracket \Delta \rrbracket, \llbracket \Gamma \rrbracket \in \text{Poset}_0$
 $\llbracket \cdot \rrbracket = 1$
 $\llbracket \Delta, x : A \rrbracket = \llbracket \Delta \rrbracket \times \llbracket A \rrbracket$
 $\llbracket \Gamma, x : A \rrbracket = \llbracket \Gamma \rrbracket \times \llbracket A \rrbracket$

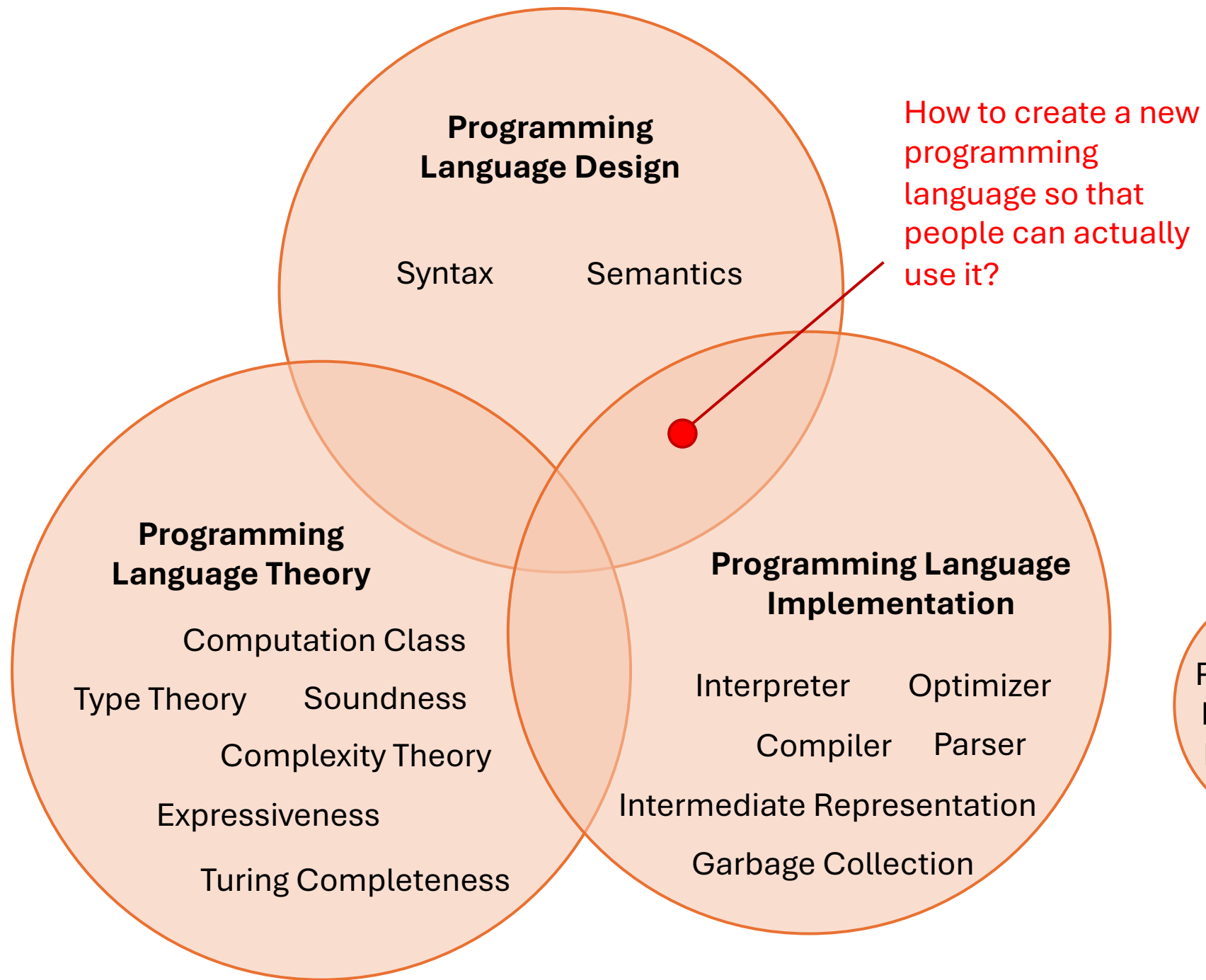
Copies from register $\overline{s_n}$, truncating if the destination is smaller.
 Lexicographically sorts the table with columns $\overline{b_n}$ and $\overline{a_n}$.
 Merges adjacent duplicate rows via σ from the unique elements s .
 Merges two sorted tables with columns $\overline{a_n}$ and $\overline{b_n}$.

Lobster

Scallop



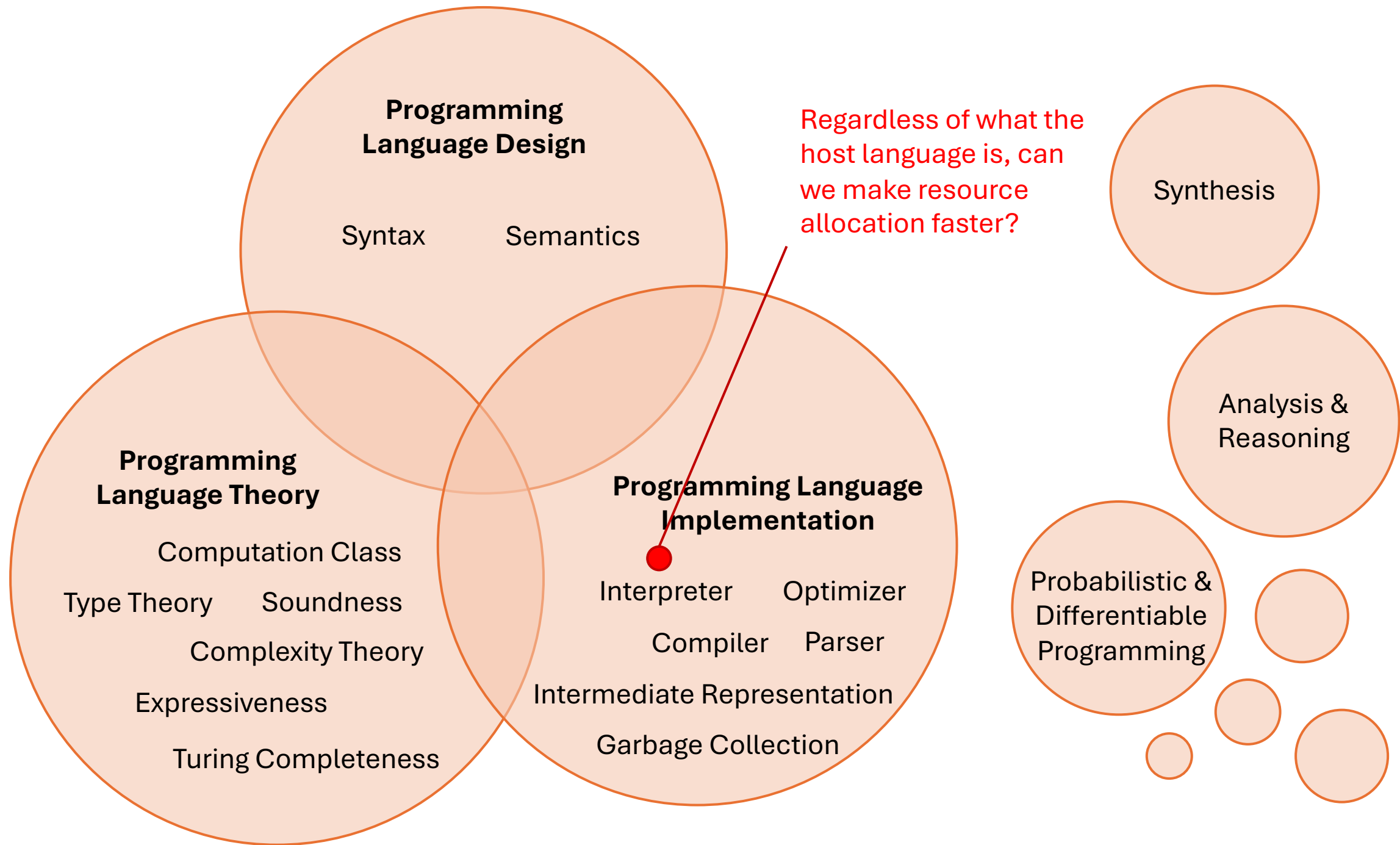




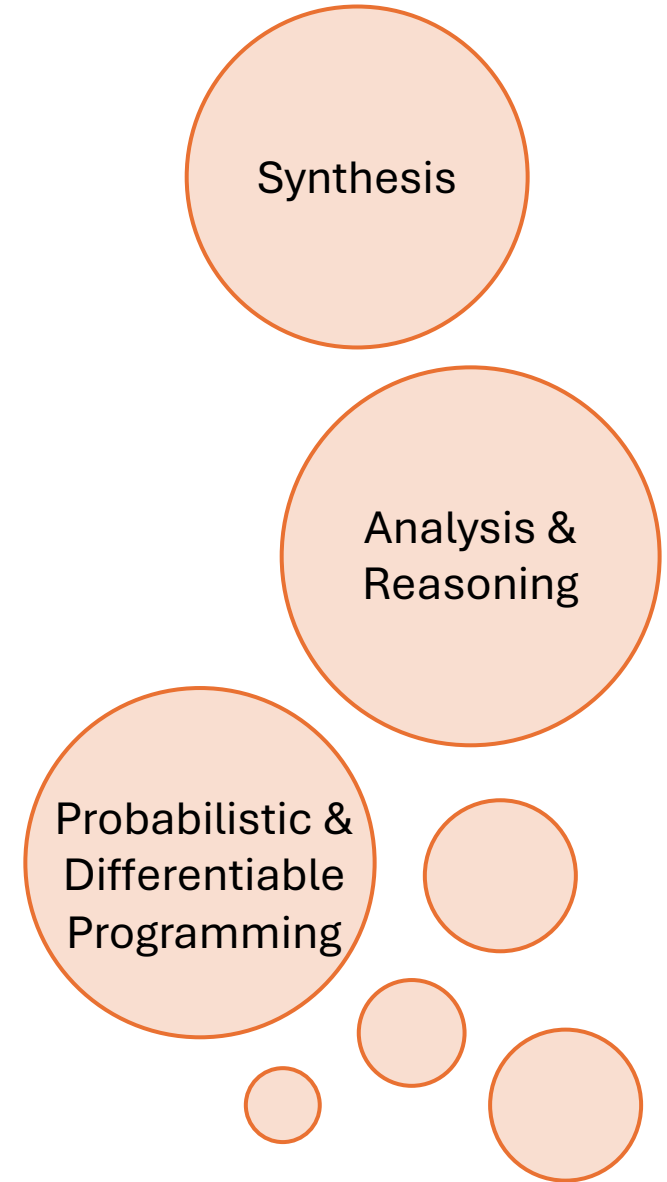
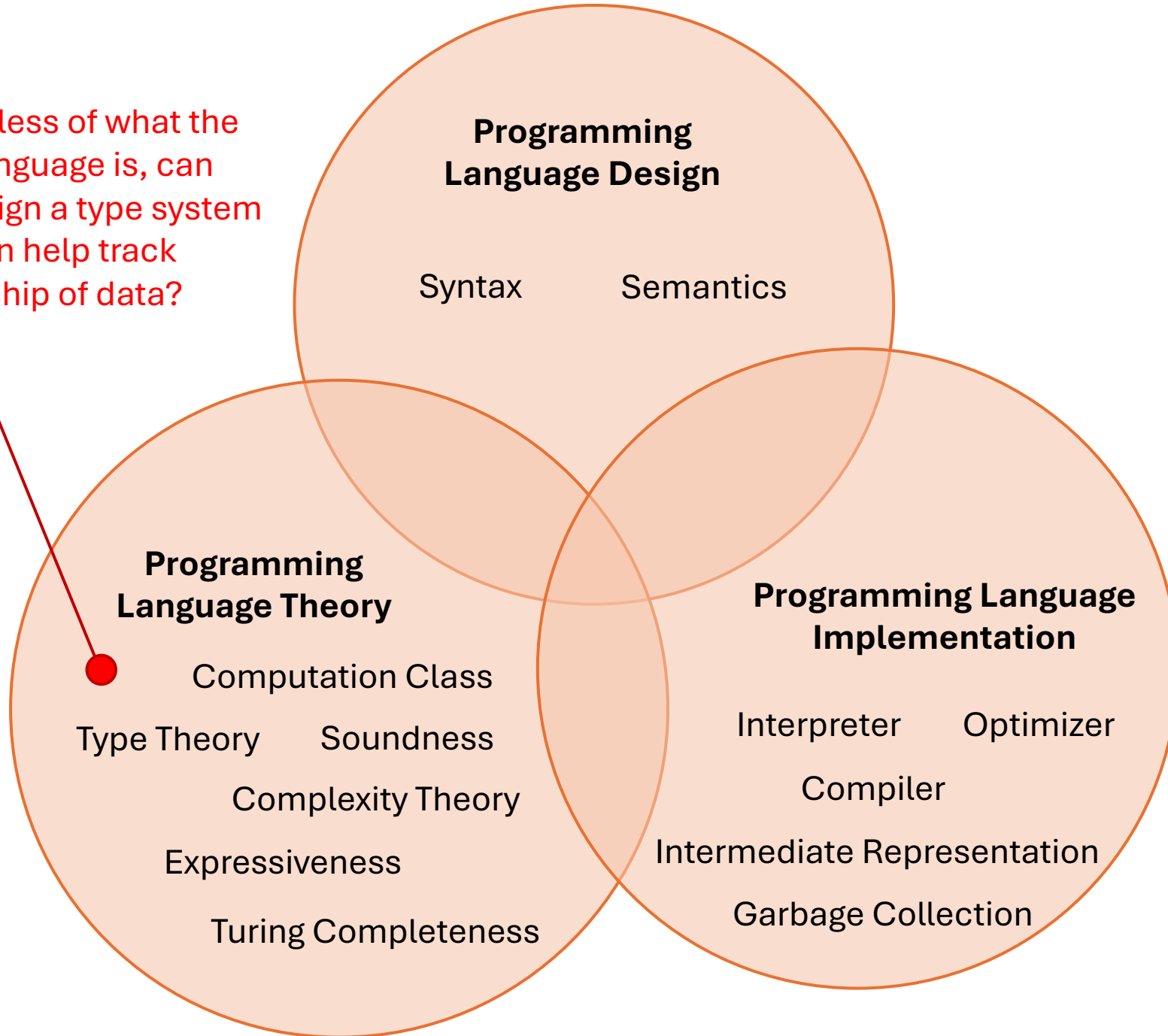
Synthesis

Analysis & Reasoning

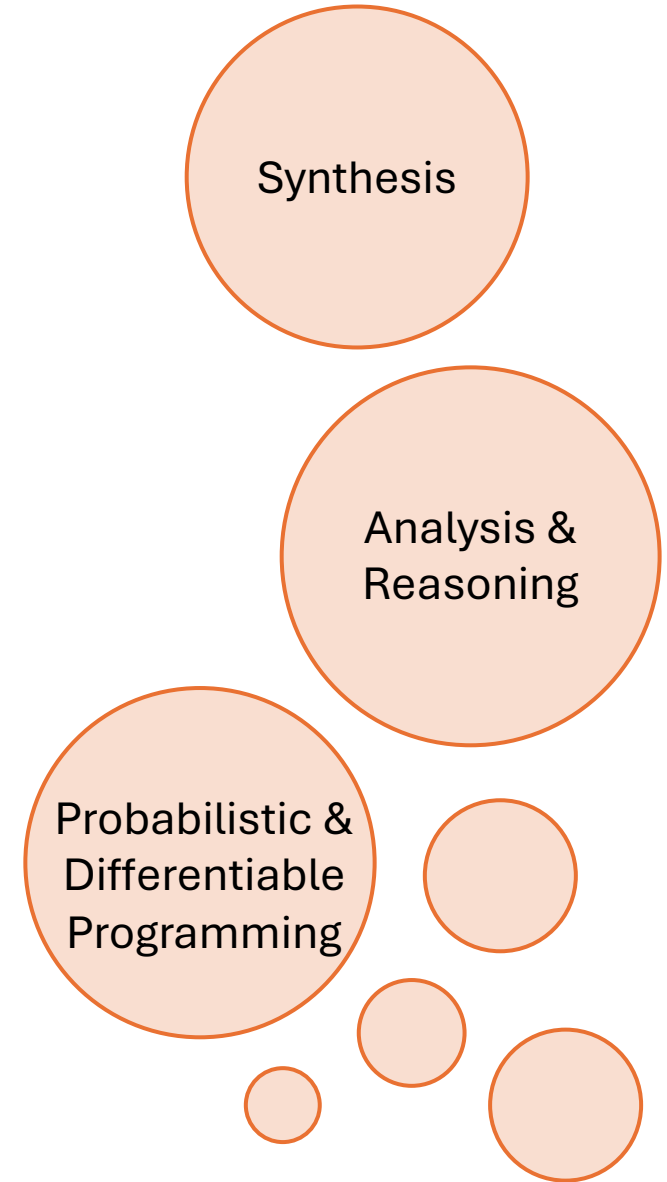
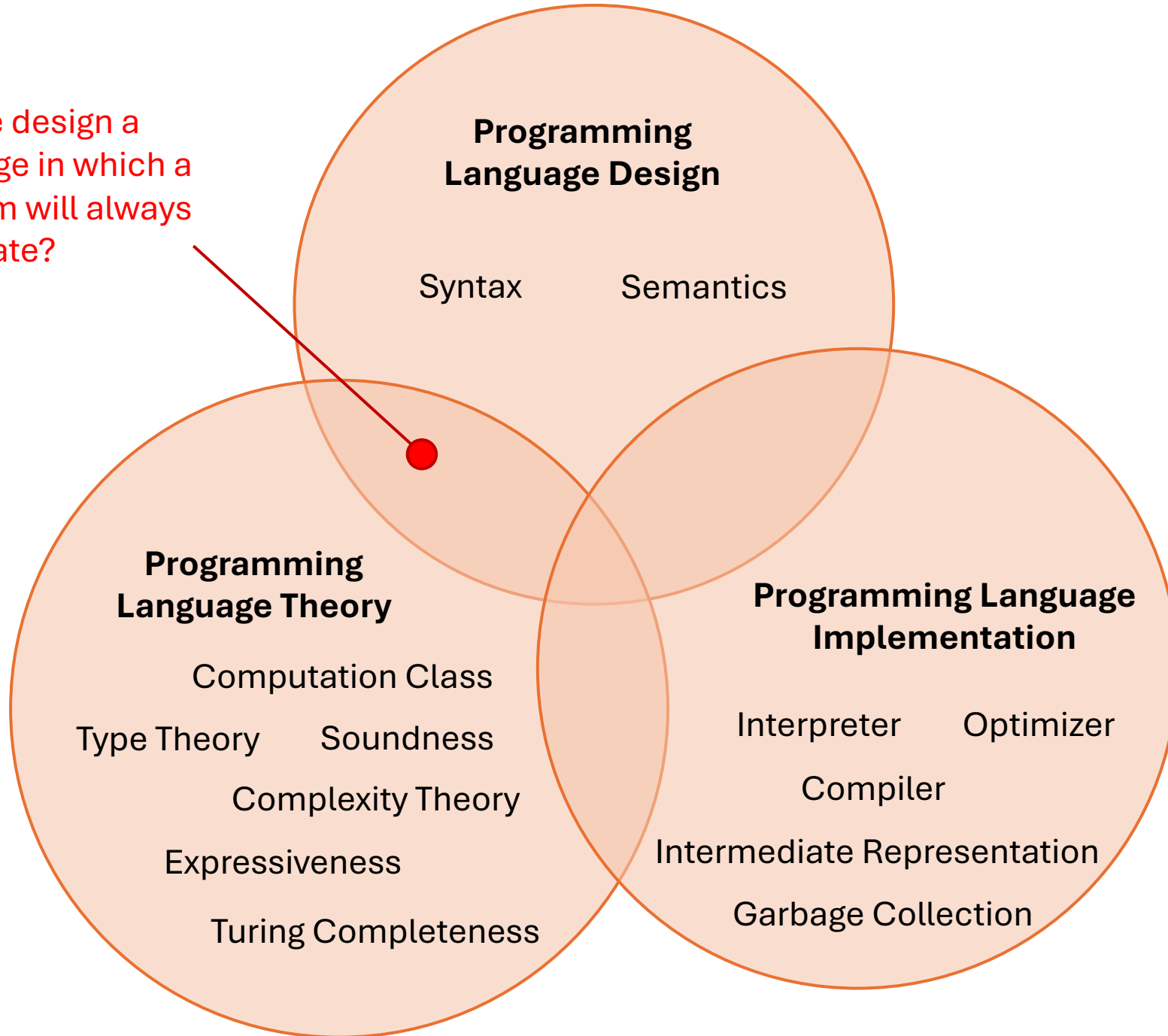
Probabilistic & Differentiable Programming

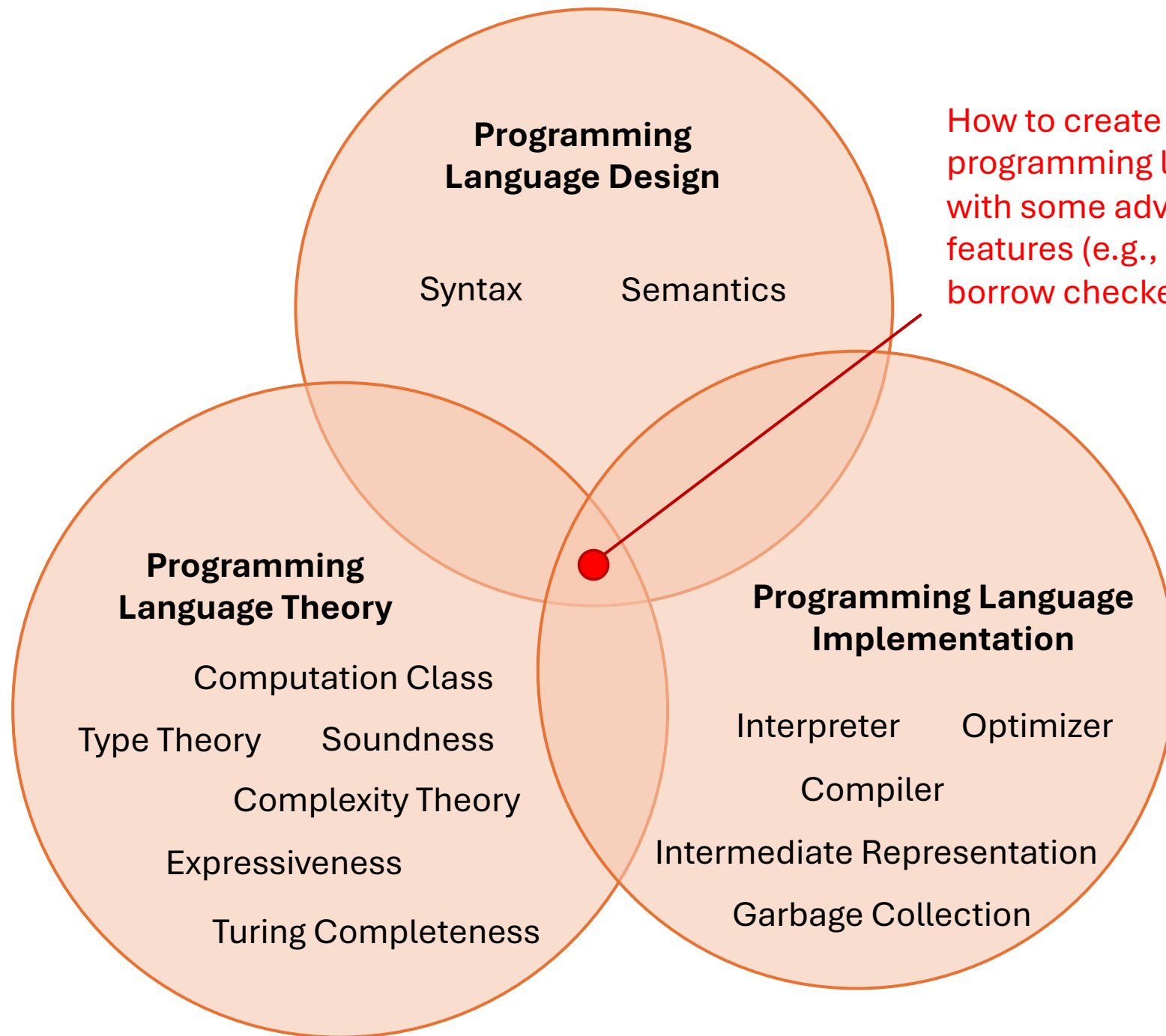


Regardless of what the host language is, can we design a type system that can help track ownership of data?



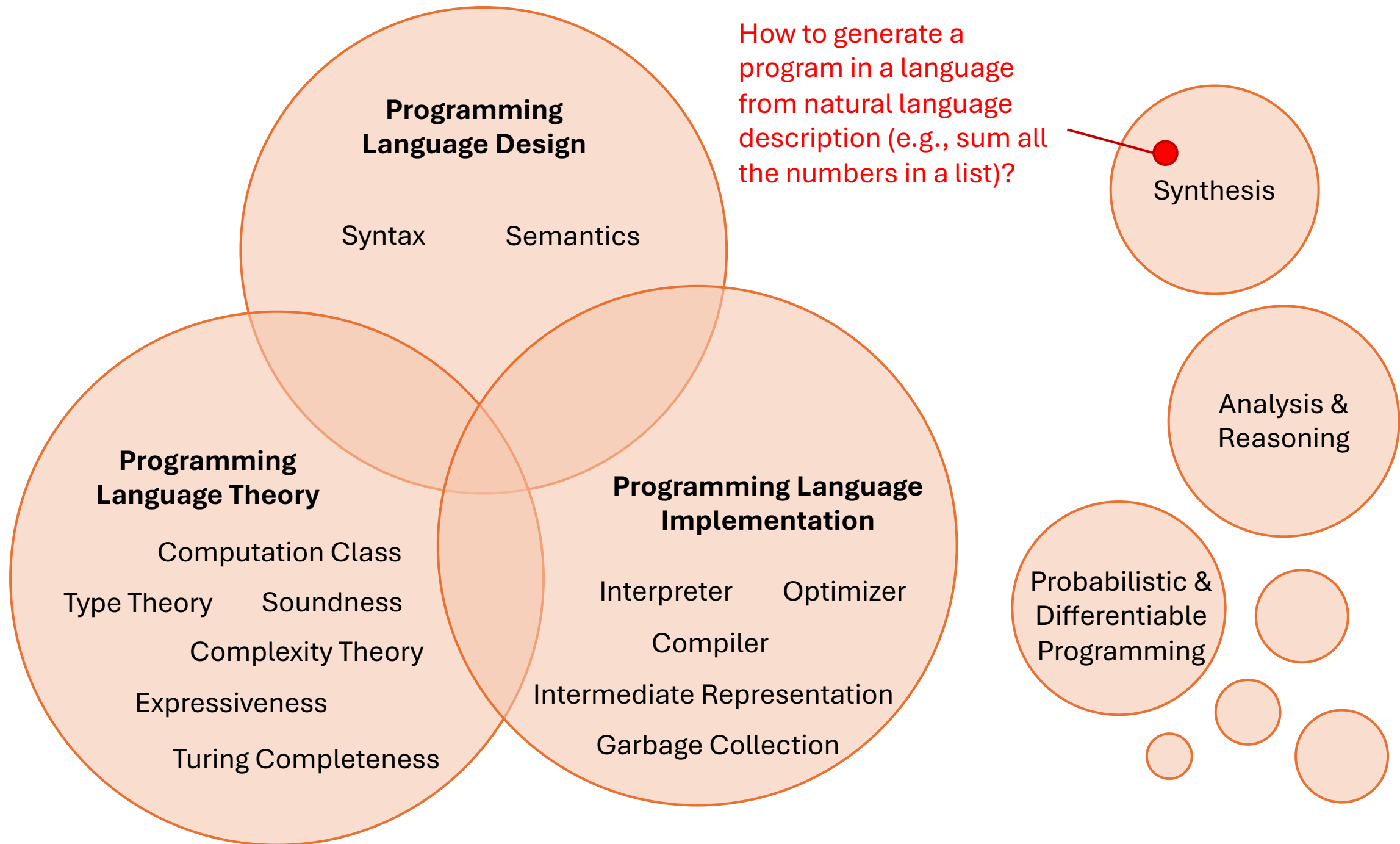
Can we design a language in which a program will always terminate?

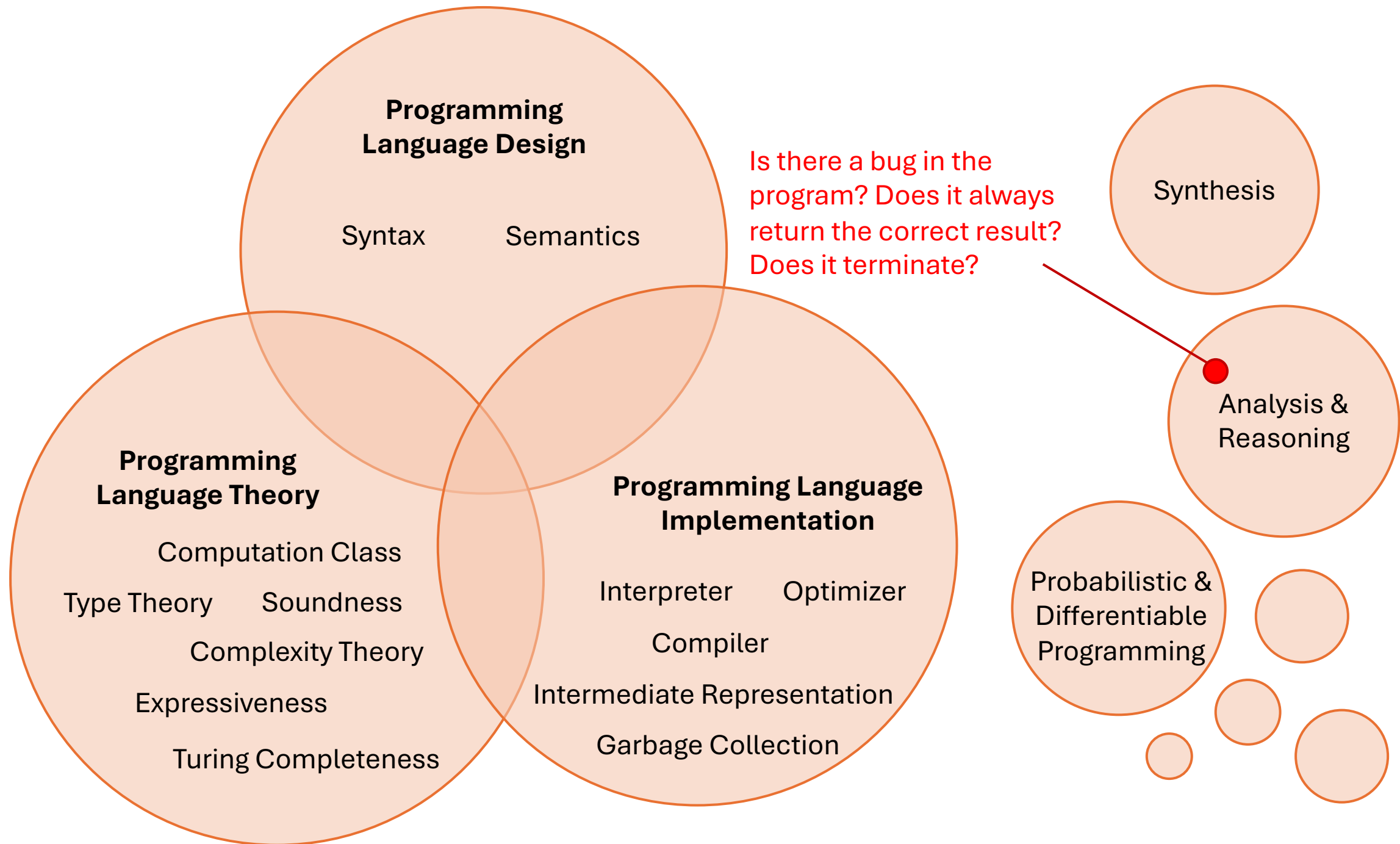


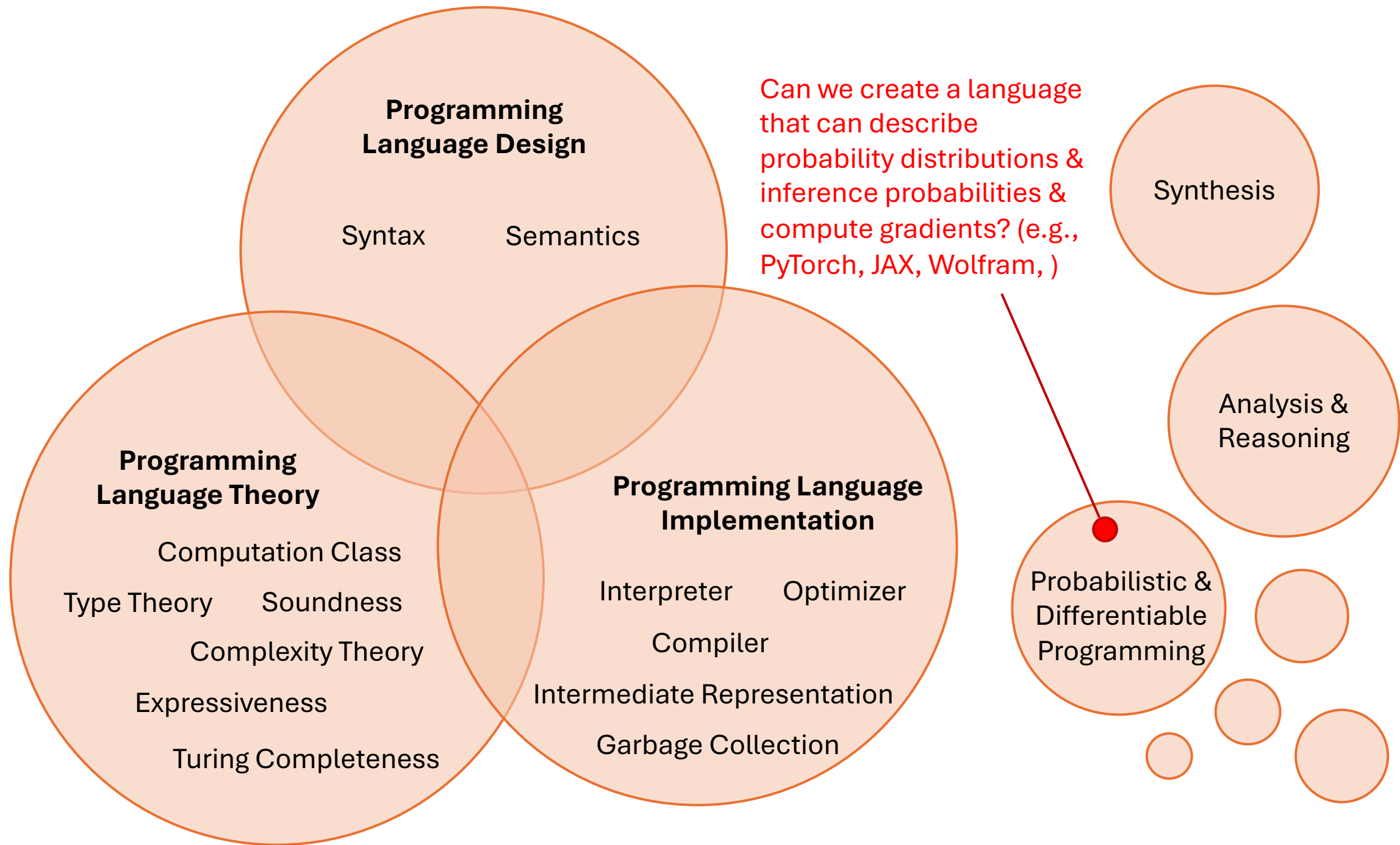


How to create a new programming language with some advanced features (e.g., Rust's borrow checker)?









Logistics

- Lecture
 - When: Tue/Thu 3:00 – 4:15pm
 - Where: Krieger Hall 170
- Course Website
 - <https://pl.cs.jhu.edu/pl/>
 - Discussions: [counselore](#)
 - Grading: [gradescope](#)

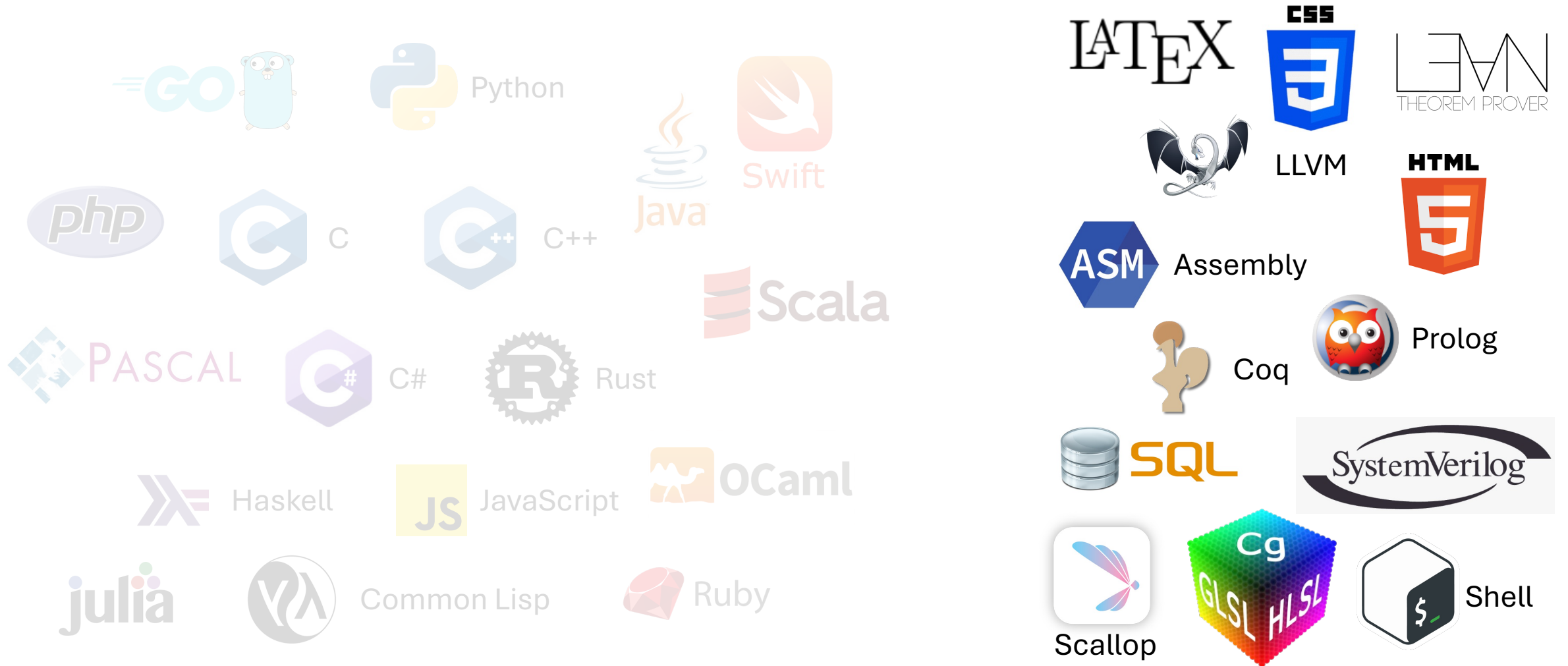
Programming Languages



General Purpose Programming Languages



Domain Specific Programming Languages



Programming Languages **that we will learn**

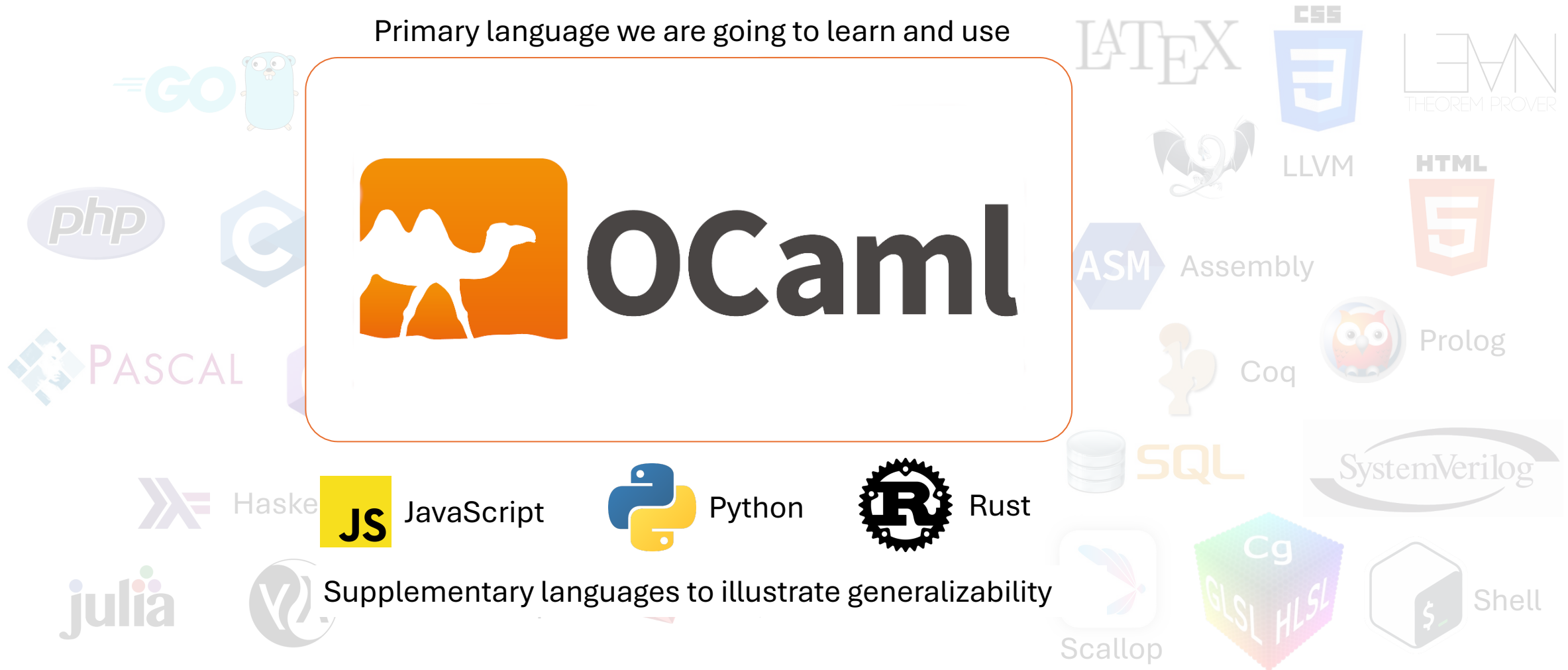


Programming Languages **that we will learn**

Primary language we are going to learn and use



Supplementary languages to illustrate generalizability



Key Learning Objectives

- **Programming Paradigms:**
 - Imperative, Functional, Declarative, ...
- **Programming Principles:**
 - Lambda Functions, Recursion, Types, Records, Algebraic Data Structures, State, Monad, Effect, Concurrency, Exceptions
- **Mathematical Foundation:**
 - Syntax, Semantics, Proof Tree, Deductive Reasoning, Equivalence
- **Programming Languages:**
 - OCaml (main), Rust, Python, JavaScript/TypeScript

What (I hope) you can do **AFTER** the course:

- **Spend less time on learning individual programming language!**
 - Just jump right in and be proficient!
 - Why? Because you already know all the principles
- **Create your own programming language!**
 - Design and implement a new language
 - Instead of making features as boilerplate APIs, turn features into a language!
- **Be better at commanding LLMs to help you write programs!**
 - Understand which language best suit your need
 - Be more precise at describing your desired program!

Grading

- Quiz & Attendance (10%)
- Homework Assignments (8 in total, 30%)
- Midterm Exam (25%)
- Final Exam (35%)

Quiz & Attendance (10%)

- 10 quizzes in total, 1% each
- Graded by completion, NOT by correctness
- The quiz will be on paper; please bring your own pen
- Quiz will be notified 1 week in advance, usually on Thursday
- Maximum 1 quiz per week
- You can miss 2 quizzes, meaning that you'll get 2% overall grade for free
- (no quiz in the first week; starting from the second week)

Quiz & Attendance (10%)

Example Quiz:

Name: _____

Complete the Proof Tree Below:

True Or (True And False) \Rightarrow False

Homework Assignments (30%)

- Throughout the semester, we will be implementing an interpreter for a **functional language** called **Fb**
 - Pronunciation: **F-Flat** (if you know music theory... this is just E)
 - Written in ASC-II characters, we use **Fb**
- Each homework assignment corresponds to one language feature within **Fb**
 - Starting from operators to state, objects, and types
- You will be able to see all the assignments from the beginning
 - Roughly speaking, you'll have 1-1.5 weeks per assignment

Homework Assignments (30%)

- The assignments are graded on GradeScope:
 - Autograders are available for most assignments; manual grading for a few
 - For autograded ones, you'll be able to see your score immediately after you submit. The score is based on the number of test cases passed
- Late Policy:
 - You will have 12 late days in total throughout the semester
 - For each assignment, you can use up to 3 late days
 - Please count the late days on your own; if the usage is exceeded, we will put 0 on some of your assignments

Homework Assignments (30%)

- **Collaboration Policy:**

- Strongly encouraged
- In the submission, please add “acknowledgement.md” when you have collaborated with other students; note their name down

- **Use-of-AI Policy:**

- Okay to use but not entirely encouraged
- In the submission, please add “acknowledgement.md” when you have used AI (LLMs, VLMs, etc.) Carefully note down the LLM and its version, copy-paste the entire conversation history (when available)

Homework Assignments (30%)

- Extra Credit

- For programming assignment 4/5/6/7/8, +1% overall grade if:
 - You have implemented the entirety of the assignment with another programming language, including runnable test cases
 - Include a simple README.md telling us how to build & run
 - Recommended language: Rust, JavaScript/TypeScript, Haskell, Python, Julia
 - Other languages that you can try working with: C/C++, Java, Scala, Go, Ruby
- If you have done all, they are +5% overall grade
- Submission using .zip by email to the instructor (please CC TA & CA)
- Submit anytime within 1 week after the corresponding assignment's due date
 - E.g., assignment 4 due Feb 18, the extra credit can be submitted before Feb 25.
- Use of LLM welcomed, please acknowledge in the submission

Midterm (25%) & Final (35%)

- Midterm: Mar 10 (Tues) in-class, 1-hour
- Final: TBD, 2.5-hour
- Study materials & practice exams will be made available online
 - Answers will not be available, but we welcome you to come to the office hour or send private posts to check with the instructors & TA & CA
- You can bring one piece of cheat sheet (size TBD)