

$$\lambda x. ((\lambda y. f y) x)$$

$$\rightarrow_{\beta} \lambda x. (\underline{f x})$$

$$\rightarrow_{\eta} f$$

Lecture 12 Encodings & Rewrites

$$e \cong e'$$

\mathbb{P}_y $x = (1, 2)$ # binary tuple
 $x = ()$ # unit
 $x = (1, 2, \text{True})$ # triplet

\mathbb{F}_b $e ::= x \mid \bar{i} \mid b \mid \dots \mid \text{Let } x = e \text{ In } e'$
 $\mid \dots$
 $\mid \underline{(e, e)}$ ① tuple constructor
 $\mid \underline{\text{Left } e} \mid \underline{\text{Right } e}$
 ② Left acc ③ Right Acc

$$v ::= \dots \mid (v, v) \quad e \Rightarrow v$$

$$[\text{Tuple}] \frac{e_1 \Rightarrow v_1 \quad e_2 \Rightarrow v_2}{(e_1, e_2) \Rightarrow (v_1, v_2)}$$

Introduced new
syntax

$$[\text{Left}] \frac{e \Rightarrow (v_1, v_2)}{\text{Left } e \Rightarrow v_1}$$

Introduced new
semantics

$$[\text{Right}] \frac{e \Rightarrow (v_1, v_2)}{\text{Right } e \Rightarrow v_2}$$

Tuples are primitive

Rewrites / Syntax Sugar

Py for x in list:

...

Java for (int i = 0; i < list.length(); i++) {

x = list[i];

} ...

$$\lambda\text{-calc} \quad \text{PAIR} \equiv_{\delta} \lambda x. \lambda y. \lambda f. f x y$$

$$\text{LEFT} \equiv_{\delta} \lambda p. p (\lambda x. \lambda y. x)$$

$$\text{RIGHT} \equiv_{\delta}$$

$$F^b \quad (e_1, e_2) \equiv_{\delta} (\text{Pair } e_1) e_2 \quad \leftarrow \text{Desugaring}$$

$$\text{Pair} \equiv_{\delta} (\text{Fun } x \rightarrow \text{Fun } y \rightarrow \text{Fun } f \rightarrow f x y)$$

$$\text{Left } e \equiv_{\delta} (\text{Left } e)$$

$$\text{Left} \equiv_{\delta} (\text{Fun } p \rightarrow p (\text{Fun } x \rightarrow \text{Fun } y \rightarrow x))$$

$$\underbrace{(e_1, e_2)}_{F^b_{\tau}} \cong \text{Pair } e_1, e_2 \quad \left(\text{where Pair is defined above} \right)$$

Definition of \cong

$$\begin{array}{ccc} (v_1, v_2) & \cong & \text{Fun } f \rightarrow f v_1 v_2 \\ \uparrow \cap & & \uparrow \cap \\ F^b_{\tau} & & F^b \end{array}$$

$$\begin{array}{c} \text{[Tuple]} \\ \text{pb} \\ \tau \end{array} \frac{e_1 \Rightarrow v_1 \quad e_2 \Rightarrow v_2}{(e_1, e_2) \Rightarrow (v_1, v_2)}$$

$$\begin{array}{c} \text{[Tuple]} \\ \text{pb} \end{array} \frac{e_1 \Rightarrow v_1 \quad e_2 \Rightarrow v_2}{(e_1, e_2) \equiv_{\delta} \text{Pair } e_1, e_2 \Rightarrow (v_1, v_2) \equiv_{\delta} \text{Fun } f \Rightarrow f v_1, v_2}$$

We want to prove this is correct.

☆

$$\begin{array}{c} \text{((Fun } x \dots)) e_1 \Rightarrow \text{Fun } y \rightarrow \text{Fun } f \rightarrow f v_1, y \quad e_2 \Rightarrow v_2 \end{array} \quad \begin{array}{c} (\text{Fun } f \rightarrow f e_1, y) [v_2/y] \\ \Rightarrow \text{Fun } f \rightarrow f v_1, v_2 \end{array}$$

1. LHS is a func 2. RHS is red to value 3. substitution

$$\left(\left(\left(\text{Fun } x \rightarrow \text{Fun } y \rightarrow \text{Fun } f \rightarrow f x, y \right) e_1 \right) e_2 \right) \Rightarrow \text{Fun } f \rightarrow f v_1, v_2$$

Rewrites & Hygiene

C Macro is not Hygienic

Rust Macro is Hygienic

Good

$$\rightarrow \text{PAIR} \equiv_{\delta} (\lambda x. \lambda y. \lambda f. f x y)$$

Bad

$$\text{PAIR } e_1, e_2 \equiv_{\delta} (\dots) e_1 e_2$$

$$\rightarrow \text{PAIR}' e_1, e_2 \equiv_{\delta} \lambda f. f (e_1) (e_2)$$

$$e_1 \equiv f$$

$$\text{PAIR } f e_2 \equiv_{\delta} (\lambda x. \lambda y. \lambda f. f x y) f e_2$$

$$\equiv_{\alpha} (\lambda x. \lambda y. \lambda g. g x y) f e_2$$

$$\equiv_{\beta} (\lambda y. \lambda g. g f y) e_2$$

$$\text{PAIR}' f e_2 \equiv_{\delta} \lambda f. f f e_2$$

$$\equiv_{\delta} (\lambda t. \boxed{\text{IF TRUE}} (\text{LEFT } t) (\text{RIGHT } t)) (\text{PAIR } a b) \Rightarrow a$$

$$\equiv_{\beta} (\lambda b. \lambda t. \lambda f. b t f) (\lambda x. \lambda y. x)$$

$$\equiv_{\beta} \lambda t. \lambda f. ((\lambda x. \lambda y. x) t f)$$

$$\equiv_{\beta} \dots \equiv_{\beta} \boxed{\lambda t. \lambda f. t}$$

$$\equiv_{\beta} \dots \equiv_{\beta} (\lambda t. \lambda f. t) (\text{LEFT } t) (\text{RIGHT } t)$$

$$(\lambda t. \text{LEFT } t) (\text{PAIR } a b)$$

$$\equiv_{\delta} (\lambda t. \text{LEFT } t) ((\lambda x. \lambda y. \lambda f. f x y) a b)$$

$$\equiv_{\beta} \dots \equiv_{\beta} (\lambda t. \text{LEFT } t) (\lambda f. f a b)$$

$$\equiv_{\eta} \text{LEFT } (\lambda f. f a b)$$

$$\equiv_{\delta} (\lambda p. p (\lambda x. \lambda y. x)) (\lambda f. f a b)$$

$$\equiv_{\beta} (\lambda f. f a b) (\lambda x. \lambda y. x)$$

$$\equiv_{\beta} (\lambda x. \lambda y. x) a b$$

$$\equiv_{\beta} \equiv_{\beta} a$$