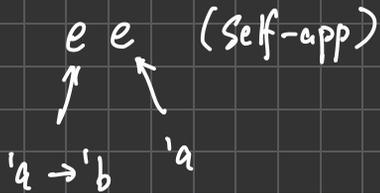


$e ::= \dots$

| Let Rec f $x = e_1 \text{ In } e_2$



fact $\stackrel{\text{def}}{=} \lambda n. n \leq 1 ? 1 : \text{fact}(n-1) * n$

fact $\stackrel{\text{def}}{=} \lambda f. \lambda n. n \leq 1 ? 1 : f(n-1) * n$

F(f) $= \lambda n. n \leq 1 ? 1 : f(n-1) * n$

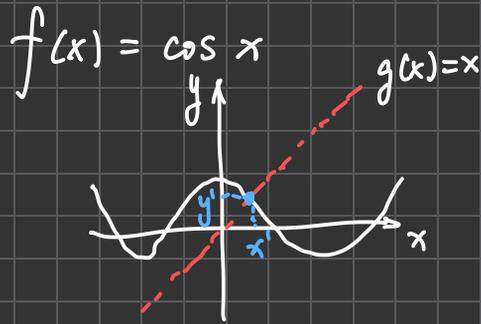
factorial : $\mathbb{N} \rightarrow \mathbb{N}$

f : $\mathbb{N} \rightarrow \mathbb{N}$

$F(f)$: $\mathbb{N} \rightarrow \mathbb{N}$

factorial = $F(f)$ = f Fix-point

if $f(x) = x$
then x is a fixpoint of f



x' is a fix point of f

$$\underset{\Delta}{\text{Fixpoint}}(f) = x$$

$$\underset{\sim}{F}(f) = \lambda n. n \leq 1 ? 1 : \underset{\sim}{f}(n-1) * n$$

$$\underset{\Delta}{F}(f) = f$$

$$(\dots) f = f$$

(Hypo) $\text{Fix} : (A \rightarrow A) \rightarrow A$

$$\text{recursive } f = \text{Fix} \left(\lambda f. \right. \\ \left. \underset{\Delta}{\lambda} \text{input}. \text{ body } \text{--} (\text{input}, f) \right)$$

$e ::= \text{Let Rec } f \ x = e_1 \ \Sigma_n \ e_2 \mid \dots \quad (F^b)$

$$\underbrace{\text{Fun } x \rightarrow \text{Let Rec } f \ x = e_1 \ \Sigma_n \ f \ x}_{(f)} / f \Rightarrow v_2$$

$$\text{[Let Rec]} \quad e_2 \left[\text{fun } x \rightarrow e_1 \left[\frac{\text{Let Rec } f \ x = e_1 \ \Sigma_n \ e_2 \Rightarrow v_2}{\text{Let Rec } f \ x = e_1 \ \Sigma_n \ e_2 \Rightarrow v_2} \right] / f \right] \Rightarrow v_2$$

Annotations for the reduction:

- Red arrows: "string / func name" pointing to e_1 , "string var name" pointing to x .
- Blue arrows: "func body" pointing to e_2 , "Continuation" pointing to f .

$$\text{[Let]} \quad \frac{e_1 \Rightarrow v_1 \quad e_2 [v_1 / x] \Rightarrow v_2}{\text{Let } x = e_1 \ \Sigma_n \ e_2 \Rightarrow v_2}$$

$$\left[\begin{array}{l} \text{Let Rec odd } x = \\ \left\{ \begin{array}{l} \text{If } \boxed{x} = 1 \ \text{Then True} \\ \text{Else If } \boxed{x} = 0 \ \text{Then False} \\ \text{Else odd } (\boxed{x} - 2) \end{array} \right. \Rightarrow \text{False} \\ \Sigma_n \\ \text{Odd } 2 \end{array} \right.$$

Evaluation trace:

- False
- ↑
- Odd 0
- ↑
- Odd 2

$$\frac{\text{false}}{e_1 [0 / x]} \Rightarrow \text{false}$$

⋮

$$\text{(odd 0)} [\dots / \text{odd}]$$

$$\frac{(\text{Fun } x \rightarrow \dots) \Rightarrow \star \quad 2 \Rightarrow 2 \quad e_1 [\dots \text{ In odd } x / \text{odd}]}{\Gamma 2 / x] \Rightarrow \text{false}}$$

$$\text{Fun } x \rightarrow \text{Let Rec odd } x = e_1$$

$$(\text{Fun } x \rightarrow e_1 [\text{In odd } x / \text{odd}]) 2 \Rightarrow \text{false}$$

⋮

$$\text{Fun } x \rightarrow \text{Let Rec odd } x = e_1$$

$$\frac{(\text{odd } 2) [\text{Fun } x \rightarrow e_1 [\text{In odd } x / \text{odd}] / \text{odd}]}{\Rightarrow \text{false}}$$

$$\text{Let Rec odd } x = e_1 \text{ In } \text{odd } 2 \Rightarrow \text{false}$$

Mutual Recursion

function f() {

g();

}

function g() {

f();

}

Let $f(x) = e_1$ In e_2 ← cannot have f
 ← can have f

Let Rec $f(x) = e_1$ In e_2
 ← can have f

{ Let Rec $f(x) = e_1$ In
 Let Rec $g(y) = e_2$ In
 e_3

Ok and { Let rec $f(x) = e_1$
 And $g(y) = e_2$
 In e_3

Record : { *can contain f and g*

f : fun x → e₁

g : fun y → e₂

⋮

}