A Component Security Infrastructure

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Cells @ FCS 2002

Motivation

Secure software systems

Build software systems using components





Secure software systems at the component level

Goals for a Component Security Infrastructure

- Simplicity
 - Complex protocols will be misused
- Generality
 - Applicable across a wide range of domains
- Interoperability
 - Security policies shared between components, others
- Extensibility

Evolves as component architecture evolves

Background: Cells

 A new distributed component programming language [Rinat and Smith, ECOOP2002]



Background: SDSI/SPKI

- Basis of our security infrastructure
- Features
 - Principal with public/private key pair
 - Decentralized name service
 - Extended names, name certificate
 - Group membership certificate
 - Access control
 - Principal with ACL
 - Delegation model: authorization/revocation certificate

Principles of Component Security

- Each component should be a principal
 - Traditional principals: users, locations, protection domains, ...
 - New idea: Components as principals
- Components are known to outsiders by their public key
- Components each have their own secured namespace for addressing other components
- Components may be private

Cell Identifiers: CID

- CID = the public key in the key pair generated by public key cryptosystem
 CID is a secured cell identity
- Universally unique
 - No two cells share the same CID
- Outgoing messages signed by CID-1 and verified by CID

CVM Identity



With President Cells:

 Universe is homogeneously composed of cells

Locations are also principals

- Locations are represented by cells and each cell is a principal
- Unique CVM identity via its President

Cell Header Security Information



Identity/Key Naming Lookup Table Security Policy Table Certificate Store (Delegation)

Cell Reference

Unifies many notions in one concept:

A locator of cells

Cell CID CVM CID

Network Location A capability to a cell
 No cell reference, no access

 A programming language construct: reference

 Corresponds to a SDSI/SPKI principal certificate

Name Services

- CIDs vs. Names
 - CIDs serve as universal identifiers, but names are still necessary
 - Extended name mechanism enables a cell to refer to another cell even if its CID is unknown
- Our name service is based on SDSI/SPKI
- Improvements:
 - Fewer certificates needed due to on-line nature
 - More expressive lookup algorithm

SDSI/SPKI Extended Names



SDSI/SPKI Groups





Cell Naming Lookup Table

- Online nature makes local name certificates unnecessary, unlike SDSI/SPKI
 – More suited for mobility
- Maintained by naming lookup interface, a concept closer to programming languages
- Naming entries can be effectively secured by using hooks
- Compatible with SDSI/SPKI



A More Expressive Algorithm



9/24/15



Cycle Detection Sketch

A cycle exists

Same local name expansion entry encountered twice

Solution:

Keep track of the path

 Raise an exception if the same name encountered twice

Security Policy

Each cell holds a security policy table, SPT.
Each policy is a 5-tuple.

subject	resource		access right	hook	deleg bit
	owner	unit			
Bob	thiscell	connector1	connect	NULL	0
Group1	thiscell	service1	invoke	NULL	1
Alice	Tim	service1	invoke	h	0

Subjects, Resources, Access Rights

Subjects

- Cells and a group of cells
- Local names, extended names, cell references

Resources

- Services, connectors, operations
- Partial order relations among them
- Access rights
 - Connect and invoke
- Application level protection: meaningful services and meaningful connections.

Hooks

- Designed for fine-grained access control
 - Protect a naming lookup entry: lookup("Tony")
 - Protect a specific file: read("abc.txt")
- Associated with operations
- Operation parameters verified via a predicate
- Predicate checked when the associated operation is triggered

– Example:

Hook_{lookup}(arg1) = { arg1="Tony" }



Cell Delegation

- Implements SDSI/SPKI delegation
- Each cell holds all certificates (both delegation and revocation) in a certificate store.
- Security policy table supports delegation
 - The owner of the resource might not be thiscell
 - The delegation bit indicating whether certificates can be further delegated
- Certificates are implicitly passed for delegation chain detection
 - No need for manual user intervention

Goals Revisited

Simplicity

- No complex algorithms/data structures
- Clearly defined principals and resources

Generality

- Not just cells, but components in general
- Not limited to certain applications
- Interoperability
 - Built on SDSI/SPKI standard
 - Communicate with any infrastructure that supports SDSI/SPKI
- Extensibility
 - Consideration for future additions: mobility, etc

Future Work

Security for Mobile Components

Cells can migrate
Mobile devices, PDAs

Hierarchical Security Policy
Interoperability

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Dynamic Component

 Components are named, addressable entities, running at a particular location.

 Components have interfaces which can be invoked.

 Components may be distributed across the network

Summary

- Security infrastructure in a component programming language
- Cell identity and CVM identity (president cell)
- Naming lookup table/interface
 - More expressive lookup algorithm and cycle detection
- Fine-grained access control
- Unification of security artifacts and programming language ones
- Formalization of SDSI/SPKI
- API from programming language perspective

Traditional Security Model



...Fails for Mobile Devices



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...Adapts Well with Mobile Devices



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Extended Name

An extended name is a sequence of local names $[n_1, n_2, ..., n_k]$, where each n_{i+1} is a local name defined in the name space of the cell n_i .

Example: Traditional Security Model



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...Fails in Cell Migration



Example: Cell Security Infrastructure



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...Adapts Well in Cell Migration

